

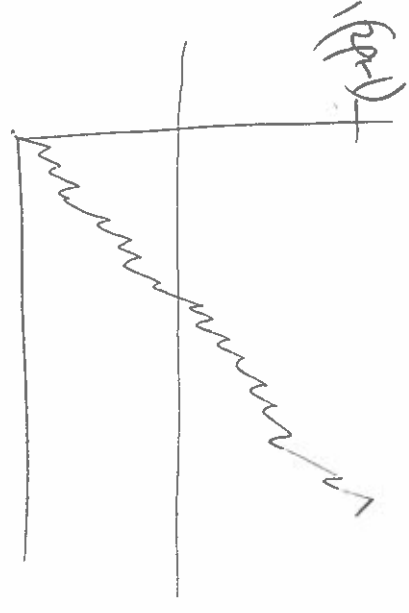
# Note on recurrence

$T =$  return time to  $0 = \min \{n > 0 : X_n = 0\}$

for a RW on  $\mathbb{Z}$ .

$$P(T = \infty) = \begin{cases} 0 & \text{no bias} \\ > 0 & \text{bias} \end{cases} \quad p \neq \frac{1}{2}$$

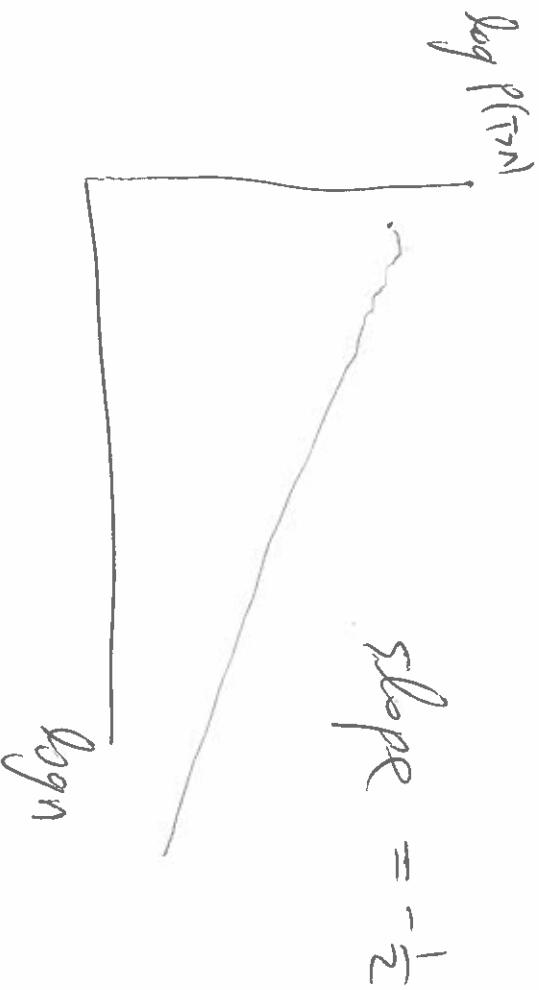
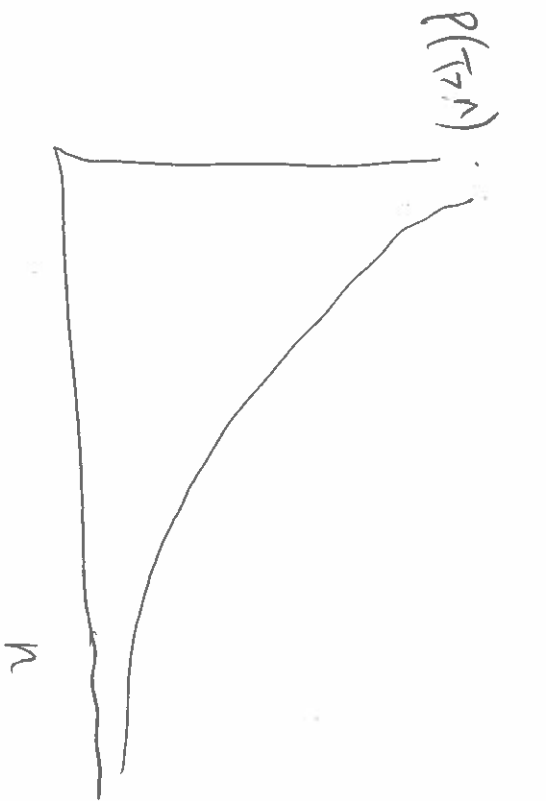
$p > \frac{1}{2}$  avg time at each level is  $\frac{1}{2p-1}$



$$\left[ \frac{1}{2p-1} = M = \frac{1}{1-u} \quad \text{so } u = 2-2p \right]$$

LLN:  $S_n \sim n(2p-1)$

If  $P = 1/2$ ,  $T < \infty$  always, but can be quite large  
 $P(T > n) \approx \frac{c}{\sqrt{n}}$   $E T = \infty$



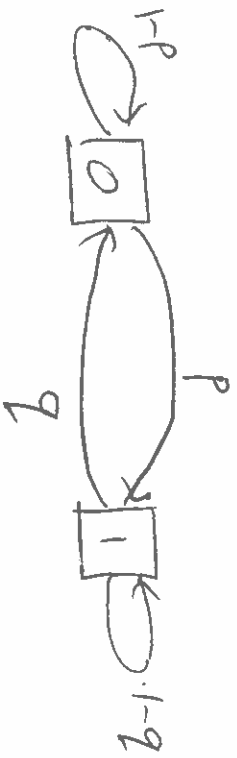
$$\log P(T > n) \approx c' + \frac{1}{2} \log n$$

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Note:  $\ln z^2 \quad P(T > n) \approx \frac{c}{\log n}$

# Stationary distributions

eg. machine is working (1) or broken (0).



$$P = \begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix}$$

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Assume  $X_0$  has distrib.  $V^{(0)}$  ;  $P(X_0=i) = V_i^{(0)}$

What is the distrib. of  $X_n$  denoted  $V^{(n)}$  ?

$n$ -step trans probab. are  $P_{ij}^n = P(X_n=j | X_0=i)$

$$\text{So } P(X_n=j) = \sum_i P(X_0=i) \cdot P(X_n=j | X_0=i)$$

$$= \sum_i V_i^{(0)} P_{ij}^n = (V^{(0)} P^n)_j$$

Summary:  $V^{(n)} = V^{(0)} P^n$

$V$  is  $1 \times |S|$  vector  
 $P$  is  $|S| \times |S|$

If  $X_0$  known to be  $i$  then  $V^{(0)} = (0, \dots, 1, \dots, 0, \dots, 0)$

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e.g. 2-state M.C., with  $X_0=0$ :

$$V^{(n)} = \begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix}^n$$

e. values:  $\lambda$  s.t.  $P - \lambda I$  is singular:

$$0 = \det(P - \lambda I) = \begin{vmatrix} 1-p-\lambda & p \\ q & 1-q-\lambda \end{vmatrix}$$

$$\lambda = 1 \text{ or } \lambda = 1 - p - q$$

e. vectors:  $x$  s.t.  $xP = \lambda x$

$$\begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix} \begin{pmatrix} 1-p-q \\ 1-q \end{pmatrix} = (1-p-q, p+q-1) = (1-p-q) \begin{pmatrix} 1 & -1 \end{pmatrix}$$

$$\begin{pmatrix} q & p \\ q & 1-q \end{pmatrix} \begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix} = (q \quad p) \quad \text{so e vectors are}$$

$$\pi = (q \quad p) \quad p = (1 \quad -1)$$

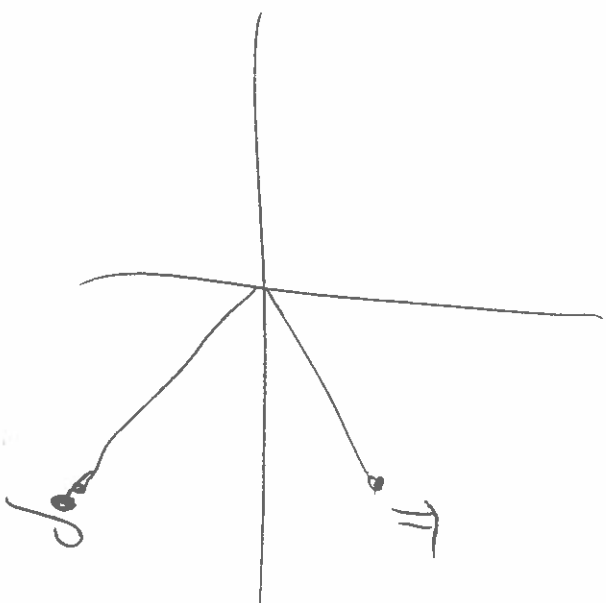
$V = a\pi + bp$  for some  $a, b$

$$V^0 P^n = a\pi P^n + b p P^n$$

$$= a\pi + b p (1-p-q)^n$$

$$\pi P^n = \pi$$

$$p P^n = p (1-p-q)^n$$



$|1-p-q| < 1$  unless  $p=q=0$  or  $p=q=1$

$$\text{So } (-p-q)^n \xrightarrow{n \rightarrow \infty} 0$$

$$\text{So } v^{(n)} \xrightarrow{n \rightarrow \infty} a\pi$$

Sum of entries is 1 so  $a = \frac{1}{p+q}$ .

$$\text{For any } v^{(0)} \text{ we have } v^{(n)} \xrightarrow{n \rightarrow \infty} \left( \frac{q}{p+q} \quad \frac{p}{p+q} \right)$$

This is the limit distrib. It gives the fraction of time at each state in the long run.

Def: A recurrent state  $i$  is positive recurrent if  $E T < \infty$   
null recurrent if  $E T = \infty$   
[ $T =$  return time]  
motivation: M.C. spends a positive fraction of time  
at a pos. recur. state

e.g. RW on  $\mathbb{Z}$  or  $\mathbb{Z}^2$  has null recurrent states.  
e.g. If  $|S| < \infty$  then only pos. recur. is possible.

Def: A M.C. is ergodic if it is ~~aperiodic~~ aperiodic  
and positive recurrent.

Thm  $\otimes$  For an irreducible ergodic M.C. there is a vector  $\pi$  s.t.  $P(X_n = j) \xrightarrow{n \rightarrow \infty} \pi_j$  for any  $V^{(j)}$

$$\left[ P_{ij}^n \xrightarrow{n \rightarrow \infty} \pi_j \quad V_{ij} \right]$$

$\otimes$  If  $N_j(T) = \#$  visits to state  $j$  by time  $T$

$$\text{then } \frac{N_j(T)}{T} \xrightarrow{T \rightarrow \infty} \pi_j$$

$\otimes$   $\pi$  is the unique solution to  $\pi P = \pi$  and  $\sum \pi_i = 1$   
ie.  $\pi_j = \sum_i \pi_i P_{ij}$

$\pi$  is called the stationary distrib. for the M.C.



Note:  $P \begin{pmatrix} | \\ | \\ | \end{pmatrix} = \begin{pmatrix} | \\ | \\ | \end{pmatrix}$  since  $P$  has row sums = 1

So  $P$  also has e-vector with  $\pi P = \pi$