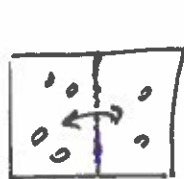


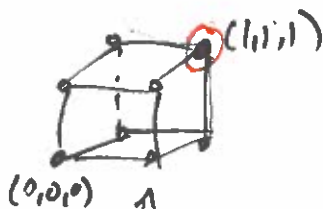
(Wednesday, March 29)

example (Ehrenfest chain)

(SRW on the hypercube) ← state space $\{0,1\}^M$



M particles



M-dimensional

$\{ (x_1, x_2, \dots, x_M) : x_i \in \{0,1\} \}$

Transitions: pick a

random index $i \in \{1, 2, \dots, M\}$, and flip the bit.

Study $X_n = \#$ particles on left side of the barrier

$$= \sum_{i=1}^M x_i \quad (x_i = 1 \text{ if particle } i \text{ on the left})$$

X_n is a MC on $\{0, 1, 2, \dots, M\}$.

Find π ; is X reversible?

Solve for π using detailed balance.

Notation: $p_{i,i+1} = \frac{M-i}{M} = \alpha_i \quad (p_{i,i-1} = 1 - \alpha_i)$

(detailed balance w/
 $i=0, j=1$)

$$\pi_0 \alpha_0 = \pi_1 (1 - \alpha_1) \Rightarrow \pi_1 = \pi_0 \cdot \frac{\alpha_0}{1 - \alpha_1}$$

$$\pi_1 \alpha_1 = \pi_2 (1 - \alpha_2) \Rightarrow \pi_2 = \pi_1 \cdot \frac{\alpha_1}{1 - \alpha_2} = \frac{\alpha_0 \alpha_1}{(1 - \alpha_1)(1 - \alpha_2)}$$

$$\pi_i = \frac{\alpha_{i-1} \alpha_{i-2} \dots \alpha_0}{(1-\alpha_i)(1-\alpha_{i-1}) \dots (1-\alpha_1)} \pi_0.$$

$$\begin{aligned} \text{Plug } \alpha_i &= \frac{M-i}{M} \Rightarrow \frac{1}{M^i} (M-i+1)(M-i+2) \dots M \\ 1-\alpha_i &= \frac{i}{M} \Rightarrow \frac{1}{M^i} i(i-1)(i-2) \dots 2 \cdot 1 \pi_0 \\ &= \binom{M}{i} \pi_0. \quad \left(= \frac{M!}{i!(M-i)!} \pi_0 \right) \end{aligned}$$

$$\begin{aligned} \text{Need } \sum_j \pi_j &= 1 \Rightarrow \pi_0 = 2^{-M}, \\ \pi_i &= 2^{-M} \binom{M}{i}. \end{aligned}$$

So $\pi \sim \text{Bin}(M, \frac{1}{2})$.

X is reversible, and π is stat.
(By thm from last class.)

Q: in general, what ^{irreducible} \wedge MCs ^{on a finite state space} are reversible?

A: Iff detailed balance eqs have a solution.

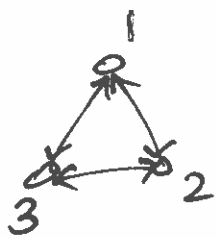
"Typically" no solutions:

$$\pi_i P_{ij} = \pi_j P_{ji} \quad - \quad n \text{ states} \\ n^2 \text{ equations}$$

Solving for $\pi = (\pi_1, \dots, \pi_n)$. n unknowns

e.g. not reversible if $p_{ij} > 0, p_{ji} = 0$
for some pair of states $i \neq j$.

e.g. $n=3$. Arb. MC on $\{1, 2, 3\}$.



$$\pi_1 p_{12} = \pi_2 p_{21} \rightarrow \pi_2 = \frac{p_{12}}{p_{21}} \pi_1$$

$$\pi_2 p_{23} = \pi_3 p_{32} \rightarrow \pi_3 = \frac{p_{23}}{p_{32}} \pi_2 = \frac{p_{12} p_{23}}{p_{21} p_{32}} \pi_1$$

$$\pi_3 p_{31} = \pi_1 p_{13}$$

$$\rightarrow \pi_1 = \frac{p_{31}}{p_{13}} \pi_3 = \frac{p_{12} p_{23} p_{31}}{p_{13} p_{32} p_{21}} \pi_1$$

$$\Rightarrow p_{12} p_{23} p_{31} = p_{13} p_{32} p_{21}$$

$$P_1(\text{one clockwise cycle}) = P_1(\text{one cc-wise cycle})$$

Theorem: if for every cycle $(v_0, v_1, v_2, \dots, v_{r-1}, v_r = v_0)$

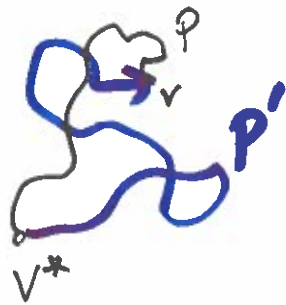
$$P(X_0 = v_0, X_1 = v_1, \dots, X_r = v_r = v_0)$$

$$\stackrel{(*)}{=} P(X_0 = v_0, X_1 = v_{r-1}, X_2 = v_{r-2}, \dots, X_{r-1} = v_1, X_r = v_0),$$

then X is time reversible.

(Converse holds: if X is reversible then $(*)$ holds for every cycle.)

Pf: (sketch) Construct π : pick a root vertex v^* . Define π_v by:



Pick path P , define

$$\pi_v = \frac{\text{probstuff}}{\text{probstuff}} \cdot \pi_{v^*}.$$

↑
"product along path P ".

Use the cycle condition to show that this definition doesn't depend on which path P you chose.

(Friday, March 31)

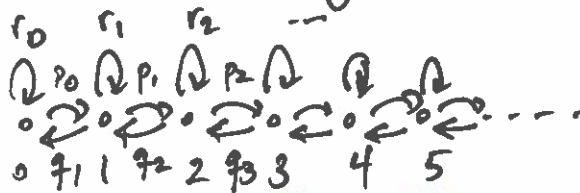
From last time: if the cycle condition^(*) is satisfied, then the MC is reversible.

$$(*) \quad P(X_1 = v_1, \dots, X_r = v_r = v_1) \\ = P(X_1 = v_1, X_2 = v_{r-1}, \dots, X_{r-1} = v_2, X_r = v_1)$$

for any cycle (v_1, v_2, \dots, v_r) .

that has a stat. prob. measure

~~Application~~ Application: any birth/death chain[^] is reversible.



$$r_i = P_{i,i} \\ p_i = P_{i,i+1} \\ q_i = P_{i,i-1}$$



cycle (2, 3, 4, 4, 3, 3, 2).

- In this case, $\pi_i = \frac{p_0 p_1 p_2 \dots p_{i-1}}{q_1 q_2 \dots q_i} \pi_0$ (directly from detailed balance) eqs. $A_0 = 1$.

~~Fact~~

Fact: if $\sum_{i \geq 0} A_i < \infty$, then chain is + recurrent
 $\left. \begin{array}{l} \\ = \infty, \end{array} \right\}$ then " is either null rec. or transient.

ex $p_i = p, q_i = q, r_i = 0$ ($\pi = \text{stat. measure}$)

Then $A_i = (p/q)^i, \pi_i = (p/q)^i \pi_0. \pi p = \pi$

$\sum A_i < \infty$ iff $p < q$.

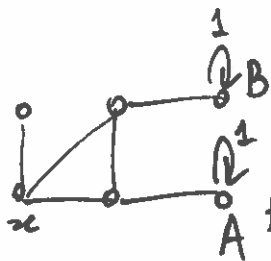
π is geometric if $p < q$.

ex $p_i = \frac{\lambda}{i+1+\lambda}, q_i = \frac{i}{i+\lambda}, r_i = \dots$ ($p_i + q_i + r_i = 1$)

Then $A_i = \frac{\lambda^i}{i!},$ so $\pi \sim \text{Poisson}(\lambda).$ $\lambda \in (0, \infty)$

(M/M/∞)
queue

Absorption probabilities / boundary value problems.



M.C. X with absorbing states A, B

$f(x) = P(X_0 = x, \text{ and } \tau_A < \tau_B),$ where

"Tau" $\tau_A = \min \{t \geq 0 : X_t \in A\}$ (same for B).
"hitting time"

(P_{xy}) transition matrix.

$$f(x) = \sum_y P_{xy} f(y), \text{ for all } x \quad + f(A) = 1$$

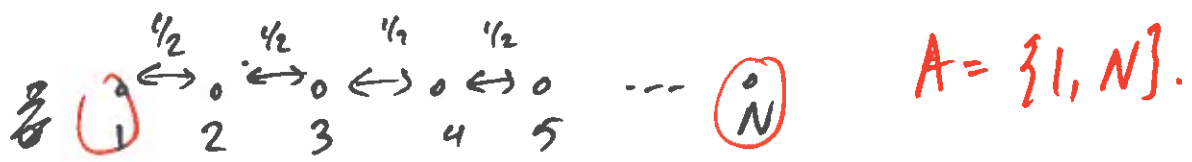
$$f(B) = 0$$

"condition on first step"

Can solve this linear system for $(f(x))_{x \in V}$.

Similar Q: what is $\mathbb{E}[T_A]$? (expected # steps before hitting A).

In general, A could be a set of states. ex SRW on $\{1, \dots, N\}$.



Similar idea: let $g(x) = \mathbb{E}_x[T_A] = \mathbb{E}[T_A | X_0 = x]$.

Similar system of eqs relating the $g(x)$'s:

$$g(x) = 1 + \sum_y p_{xy} g(y). \quad \text{for all } x.$$
$$+ g(x) = 0 \quad \text{for } x \in A.$$

$$g(x) = 1 + \frac{1}{2} g(x+1) + \frac{1}{2} g(x-1).$$

$$g(x) = x(N-x).$$

(hitting prob for $A = 1, B = N$)

$$f(x) = \frac{1}{2} f(x+1) + \frac{1}{2} f(x-1)$$