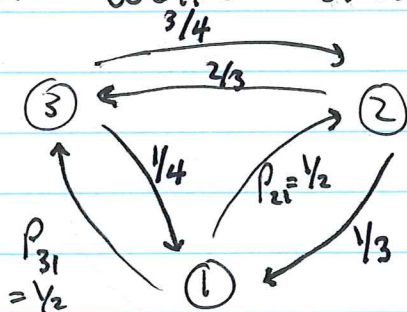


Math 152 Lecture 14

Recall: random walks: around a "state space":

Ex



$$P = \begin{bmatrix} 0 & 1/3 & 1/4 \\ 1/2 & 0 & 3/4 \\ 1/2 & 2/3 & 0 \end{bmatrix}$$

transition probabilities

$x_i^{(n)}$: probability of being in state i at time n

$\vec{x}^{(n)}$: vector of state probabilities

P_{ij} : probability of leaving state j going to state i

Model: $\vec{x}^{(n+1)} = P \vec{x}^{(n)}$

Eg. $\vec{x}^{(1)} = P \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 1/2 \\ 1/2 \end{bmatrix}$, $\vec{x}^{(2)} = P \vec{x}^{(1)} = P(P \vec{x}^{(0)}) = P^2 \vec{x}^{(0)} \approx \begin{bmatrix} 0.2917 \\ 0.375 \\ 0.3333 \end{bmatrix}$

Example from last day

(c) What is the probability of being in state 2 after a long time

$$\lim_{n \rightarrow \infty} \vec{x}^{(n)} = \lim_{n \rightarrow \infty} P^n \vec{x}^{(0)}$$

For now, numerical evidence... (math later)
(~~matlab~~ matlab/octave)

$$\vec{x}^{(n)} \rightarrow \begin{bmatrix} 0.2264 \\ 0.3962 \\ 0.3774 \end{bmatrix}$$

(d) How much does the "initial condition" $x^{(0)}$ matter? (experimentally)

$\lim_{n \rightarrow \infty} x_2^{(n)}$ seems to be 0.3962 for many different $x^{(0)}$.

"stable process"

Example notes.pdf : "specific but somewhat nerdy example" 4.1

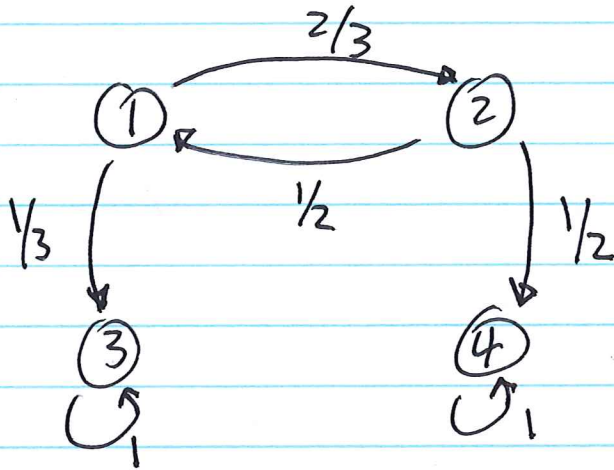
Ydnew	vs	Xavier	} see notes.pdf
<ul style="list-style-type: none"> ↳ spell works 1/2 the time. 		<ul style="list-style-type: none"> ↳ goes first ↳ spell works 1/3 of the time. 	

long duel: play until someone wins (other knocked out)
 (experimentally, each wins half the time)

short duel: Xavier as 3 attempts. after which Ydnew wins by default.

(a) describe duels as random walks.

- 4 possible states:
- ① Xavier's turn (no winner)
 - ② Ydnew's turn (" ")
 - ③ X has won
 - ④ Y has won.



$$P = \begin{bmatrix} 0 & 1/2 & 0 & 0 \\ 2/3 & 0 & 0 & 0 \\ 1/3 & 0 & 1 & 0 \\ 0 & 1/2 & 0 & 1 \end{bmatrix}$$

b) analyze probability that Y will win the short duel. (X has 3 attempts: XYXYX)

$$\vec{x}^{(0)} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$\vec{x}^{(5)} \approx P^5 \vec{x}^{(0)} \approx \begin{bmatrix} 0 \\ 7.4\% \\ 48.1\% \\ 44.4\% \end{bmatrix}$$

matlab

Y wins 52% of the time

c) long duel: matlab experiments confirm

$$\lim_{n \rightarrow \infty} P^n \vec{x}^{(0)} = \begin{bmatrix} 0 \\ 0 \\ 1/2 \\ 1/2 \end{bmatrix}$$

↑
experiments

Ex What if X and Y have a shield which blocks one ~~spell~~ successful spell ~~cost~~. cost.

States:

- ① no winner, Xavier's turn, X shield, Y shield
- ② " " " X no shield, "
- ③ " " " " Y no shield
- ④ " " " X shield, "
- ⑤-⑧ repeat but Ydnew's turn
- ⑨ Xavier has won.
- ⑩ Ydnew " " .

$P =$

0	0	0	0	1/2	0	0	0	0	0
0	0	0	0	1/2	1/2	0	0	0	0
0	0	0	0	0	0	1/2	1/2	0	0
0	0	0	0	0	0	0	1/2	0	0
2/3	0	0	0	0	0	0	0	0	0
0	2/3	0	0	0	0	0	0	0	0
0	1/3	2/3	0	0	0	0	0	0	0
1/3	0	0	2/3	0	0	0	0	0	0
0	0	1/3	1/3	0	0	0	0	1	0
0	0	0	0	0	1/2	1/2	0	0	1

X spell breaks Y shield (pointing to row 8, column 1)

X can't win on 1st round (pointing to row 8, column 1)

... matlab experiment suggests Y wins the short duel 77.8% of the time. (Xavier ~~lost~~ or didn't win)

... lim seems to have Y winning 62.5% of the time. $n \rightarrow \infty$

Transpose

Ex $A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix}$ $A^T = \begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$

def'n: if A is an $m \times n$ matrix entries a_{ij} then A^T (the transpose of A) is an $n \times m$ matrix with entries $a^T_{ij} = a_{ji}$.

- columns of A become rows of A^T .
rows " " " columns " "

- $(A^T)^T = A$ for every matrix

- save vertical space when writing
 $[1 \ 2 \ 3 \ 4]^T = \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \end{bmatrix}$

- If A is square ($n \times n$) and $A^T = A$ we say A is symmetric.

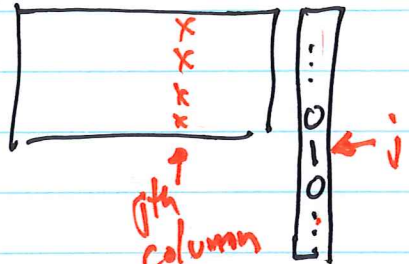
- If \vec{x} and \vec{y} are column vectors in \mathbb{R}^n then

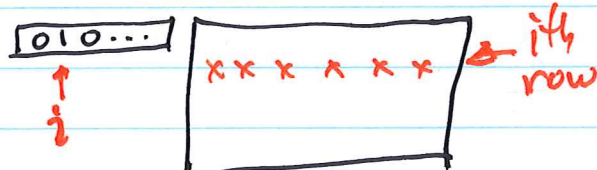
$$\vec{x}^T \vec{y} = \boxed{} = \vec{x} \cdot \vec{y} = \text{scalar}$$

- $(AB)^T = B^T A^T$ (order switches)

Operations on \vec{e}_j vectors vector of all zeros, but ~~position~~ j th position is 1.

Say A is $m \times n$ matrix.

① $A\vec{e}_j =$  $= j$ th column of A

② $\vec{e}_i^T A =$  $= i$ th row of A

③ $\vec{e}_i^T A \vec{e}_j = a_{ij}$

Note: transpose requires no arithmetic
↳ computer can even use the same memory
↳ cheap to do an ~~inverse~~ transpose.
contrast that fact with...

Matrix Inverse

Square $n \times n$ matrix A . Suppose A has rank n so $A\vec{x} = \vec{b}$ has a unique sol'n \vec{x}

(recall \rightarrow GE on Augmented matrix...)

\uparrow
told us rank

\uparrow
gave us \vec{x}

The operation of finding \vec{x} from given \vec{b} turns out to be a linear transform

Thus there is a matrix B such that

$$\vec{x} = B \vec{b} \quad (*)$$

transform

But $A\vec{x} = \vec{b}$ so multiply (*) by A :

$$A\vec{x} = AB\vec{b}$$

Therefore $AB = I$ (because $\vec{b} = AB\vec{b}$)
↑
identity matrix.

We say B is the inverse matrix of A and we write $B = A^{-1}$.