

Research Statement

Elyse Yeager

University of Illinois at Urbana-Champaign

Preliminary Examination, 08 May 2013

Disjoint Cycles and Equitable Coloring

Theorem (Corrádi-Hajnal 1963)

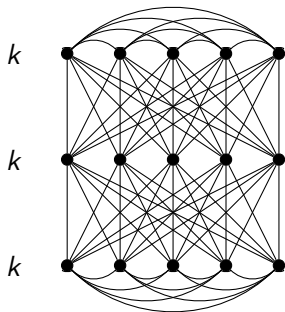
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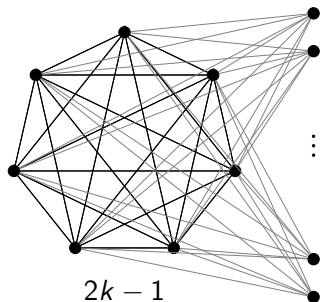
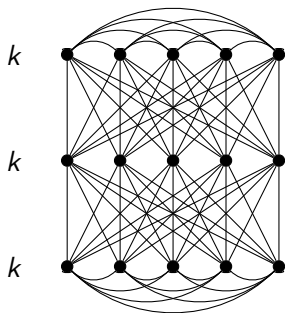


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Enomoto, Wang: Ore Version

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$$\sigma_2(G) = \min_{xy \notin E(G)} \{d(x) + d(y)\}$$

Theorem (Enomoto 1998; Wang 1999)

Let $k \geq 1$, $n \geq 3k$, and let H be an n -vertex graph with $\sigma_2(H) \geq 4k - 1$. Then H contains k vertex-disjoint cycles.

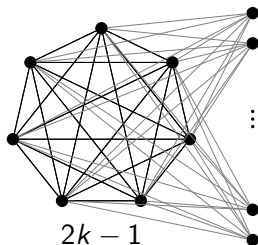
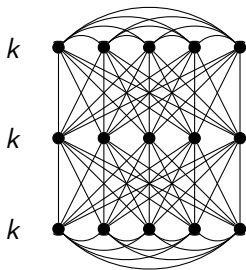
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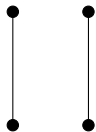
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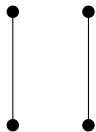


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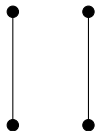
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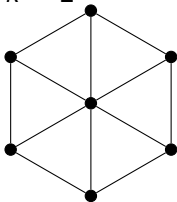
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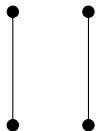


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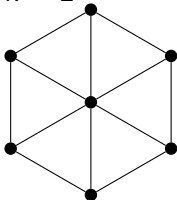
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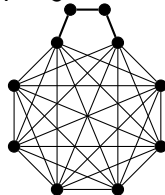
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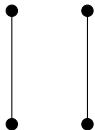


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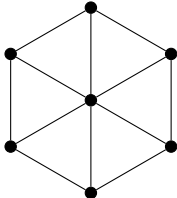
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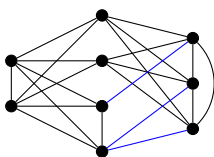
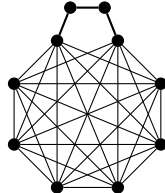
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Theorem (Kierstead, Kostochka, Y.)

Let $k \geq 3$, $n \geq 3k + 1$, and let H be an n -vertex graph with $\delta(H) \geq 2k - 1$ and $\alpha(H) \leq n - 2k$. Then H contains k vertex-disjoint cycles.

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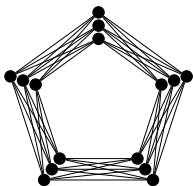
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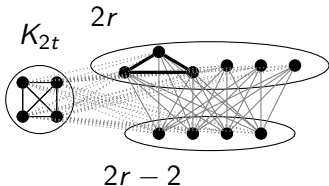
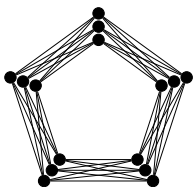
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$$n = 4(r + t) - (2t + 2)$$

Applications to Equitable Coloring

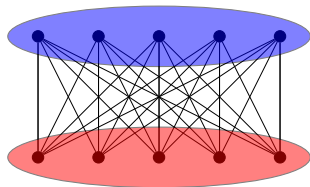
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An *equitable coloring* of a graph G is a proper vertex coloring in which the sizes of any two color classes differ by at most one.

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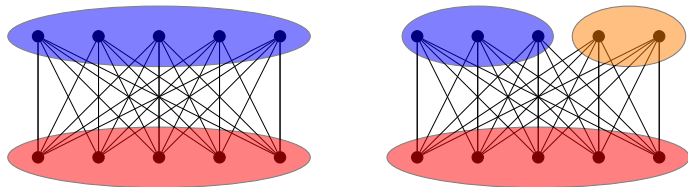
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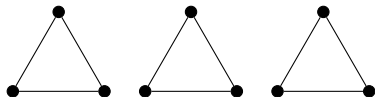
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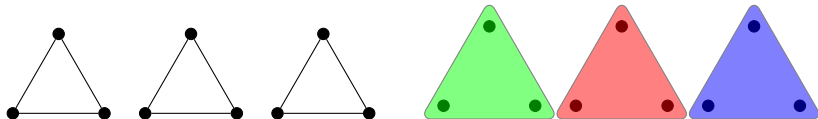


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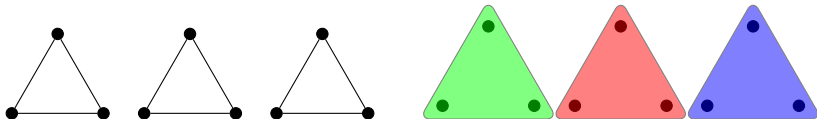


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Theorem 1 (Hajnal-Szemerédi)

If $\Delta(G) \leq k - 1$, then G is equitably k -colorable.

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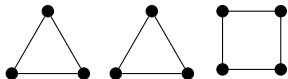
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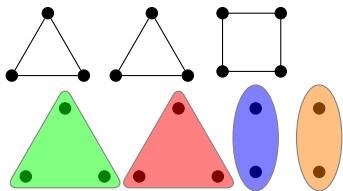
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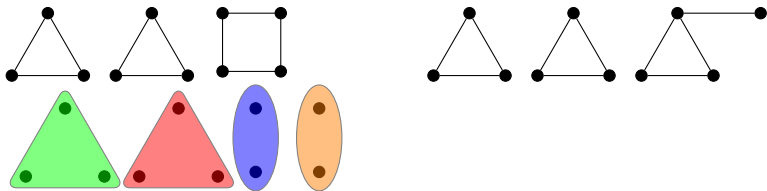
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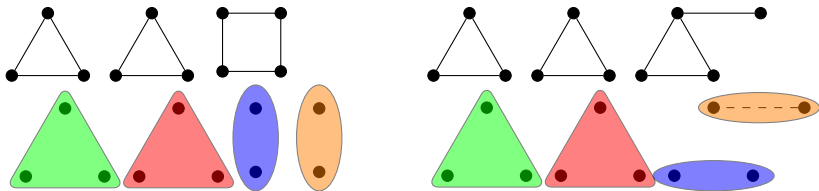
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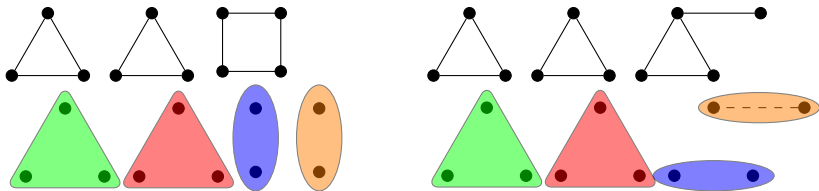
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Conjecture (Chen-Lih-Wu)

If G is a $(k + 1)$ -colorable graph with $\Delta(G) \leq k + 1$, then G has an equitable $(k + 1)$ -coloring or $k + 1$ is odd and G contains $K_{k+1, k+1}$.

Future Directions

Conjecture (Chen-Lih-Wu)

If G is a k -colorable graph with $\Delta(G) \leq k$, then G has an equitable k -coloring or k is odd and G contains $K_{k,k}$.

Refining Chen-Lih-Wu for $n = 3k$:

Conjecture

Let G be a $3k$ -vertex, k -colorable graph with $\bar{\sigma}(G) \leq 2k + 1$. If G does not contain certain subgraphs, then G has an equitable k -coloring.

Ramsey Saturation

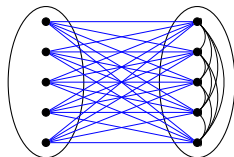
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The *saturation number* of a family of graphs \mathcal{H} , denoted $\text{sat}(n; \mathcal{H})$, is the *minimum number of edges* over all graphs G on n vertices with the property that no member of \mathcal{H} is a subgraph of G , but some member of \mathcal{H} is a subgraph of $G + e$ for every edge $e \notin E(G)$.

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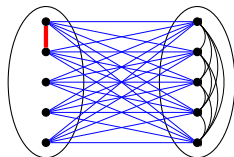
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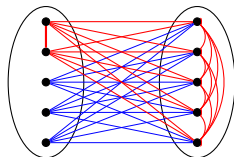
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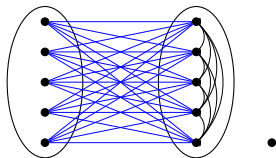
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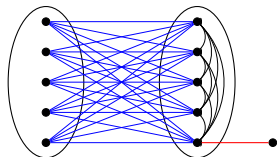
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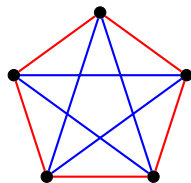
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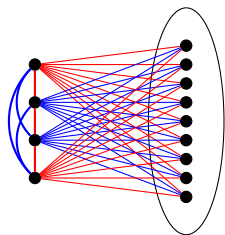
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Conjecture (Hanson, Toft)

$$\text{sat}(n; \mathcal{R}(K_{k_1}, \dots, K_{k_t})) = \begin{cases} \binom{n}{2} & n < r \\ \binom{r-2}{2} + (r-2)(n-r+2) & n \geq r \end{cases}$$

where $r := r(k_1, \dots, k_t)$.

Ramsey Saturation Number of Matchings

Theorem (FKY)

If $n > 3(m_1 + \cdots + m_t - t)$, then

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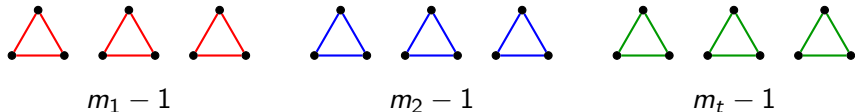
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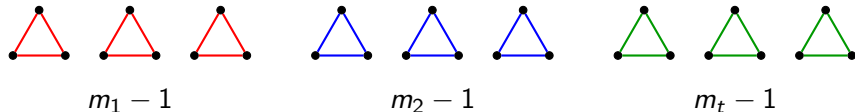
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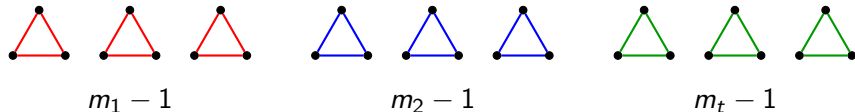
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$$r(m_1 K_2, \dots, m_t K_2) = \max\{m_1, \dots, m_t\} + 1 + (m_1 + \cdots + m_t - t)$$

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- ▶ Matchings versus other kinds of graphs
- ▶ Hanson-Toft
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 - ▶ K_4, K_4

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