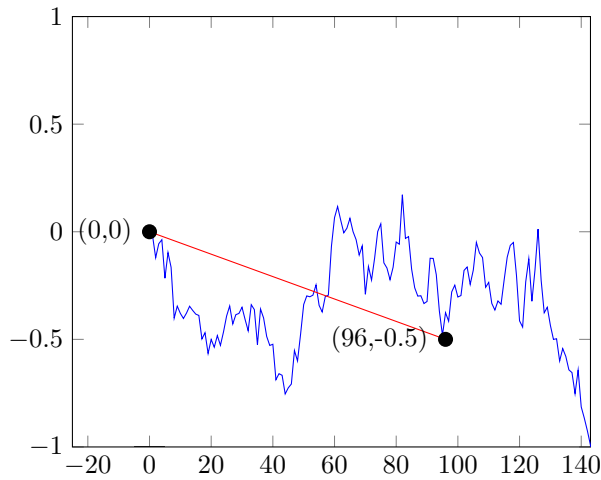


# Entropic Curves: Schrodinger Bridge

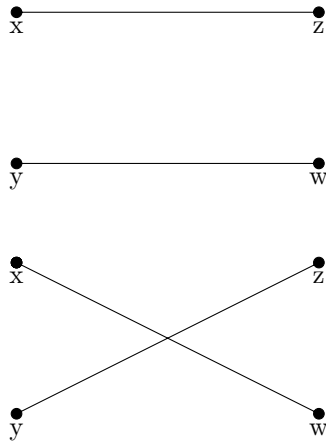
Logan Timmins

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**Figure 1.** Distance of paths connecting points  $(0,0)$  to  $(96,-0.5)$ . A **Brownian Bridge** (blue curve) and the shortest path (red).

A **Schrodinger Bridge** is a distribution of paths connecting  $x_0$  to  $x_1$  regarding more than 1 particle.



**Figure 2.** (Top) More likely paths for particles at points  $x$  and  $y$  to end at  $z$  and  $w$ , respectively. (Bottom) Potentially less likely paths for particles at points  $x$  and  $y$  to end at  $w$  and  $z$ , respectively.

Consider a collection of indistinguishable particles undergoing Brownian motion with parameter  $\varepsilon$ . We observe:

$$\begin{aligned} x_0(t=0), x_1(t=0) \dots x_n(t=0) \\ x_0(t=1), x_1(t=1) \dots x_n(t=1) \end{aligned}$$

When  $n=1$ , there's only one starting point and one end point. This distribution on paths is called a **Brownian Bridge**. When  $n$  is greater than 1, the distribution on paths is called a **Schrodinger Bridge**.

Two components of a Schrodinger Bridge

1. **Random Permutation Matrix (P)** describing which particle goes where.
2. Independent Brownian Bridges connecting each pair of particles.

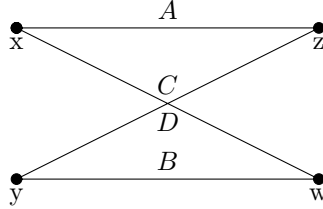
**Fact:** Solving entropy-regularized optimal transport (OT) will give coupling  $\pi_\varepsilon^*$  connecting  $\hat{\mu}_0$  to  $\hat{\mu}_1$

$$\mathbb{E}P = \pi_\varepsilon^*$$

Positions:

$$x_0(t=0), x_1(t=0) \dots x_n(t=0) \Rightarrow \hat{\mu}_0 = \sum_{i=1}^n \frac{1}{n} \delta_{x_i}(0)$$

$$x_0(t=1), x_1(t=1) \dots x_n(t=1) \Rightarrow \hat{\mu}_1 = \sum_{i=1}^n \frac{1}{n} \delta_{x_i}(1)$$



**Figure 3.** Potential paths from points  $x$  and  $y$  to  $z$  and  $w$ .

Let the path from  $x$  to  $z$  be called  $A$ , from  $y$  to  $w$  be  $B$ ,  $x$  to  $w$  be  $C$  and  $y$  to  $z$  be  $D$ . Repeat this experiment  $n$  times and record the number of times each transition ( $A, B, C, D$ ) is made. According to Brownian Motion:

$$\alpha = Prob(A) \propto e^{-\frac{\|x-z\|^2}{\varepsilon}}$$

$$\beta = Prob(B)$$

$$\gamma = Prob(C)$$

$$\delta = Prob(D)$$

$$Prob[Num(A) \approx pn, Num(B) \approx qn, Num(C) \approx rn, Num(D) \approx sn]$$

$$\approx \frac{n!}{(pn)!(qn)!(rn)!(sn)!} \alpha^{pn} \beta^{qn} \gamma^{rn} \delta^{sn}$$

By **Sterling's Formula**:  $k! \approx \exp[k \log k - k]$  Therefore,

$$\begin{aligned}
&= \exp[n \log n - n p \log n p - n q \log n q - n r \log n r - n s \log n s + n p \log \alpha + n q \log \beta + n r \log \gamma + n s \log \delta] \\
&= \exp[-n[p \log \frac{p}{\alpha} + q \log \frac{q}{\beta} + r \log \frac{r}{\gamma} + s \log \frac{s}{\delta}]] \\
&= \exp[-n KL(\begin{bmatrix} p \\ q \\ r \\ s \end{bmatrix} \mid \begin{bmatrix} \alpha \\ \beta \\ \gamma \\ \delta \end{bmatrix})]
\end{aligned}$$

This calculation tells us that when  $n$  is large, we will only see proportions  $p, q, r$  and  $s$  that are as close to  $\alpha, \beta, \gamma,$  and  $\delta$  as possible.

**Recall:**  $KL(\mathbb{P}|\mathbb{Q}) \geq 0$  or  $= 0$  iff  $\mathbb{P} = \mathbb{Q}$

**Note:**  $\begin{bmatrix} p & r \\ q & s \end{bmatrix}$  Should be a coupling of  $\hat{\mu}_0 \rightarrow \hat{\mu}_1$  and  $\begin{bmatrix} \alpha & \gamma \\ \beta & \delta \end{bmatrix}$  shouldn't be a coupling of  $\hat{\mu}_0 \rightarrow \hat{\mu}_1$  because they don't have the right marginals. Now, we'll reformulate entropic optimal transport:

$$\begin{aligned}
&\min_{\pi} \mathbb{E}_{\pi} c(x, y) - \varepsilon H(\pi) \\
&(x, y) \sim \pi, x \sim \mathbb{P}, y \sim \mathbb{Q}
\end{aligned}$$

For **finitely** supported distributions:

$$\begin{aligned}
\mathbb{P} &= \sum_{i=1}^n \frac{\delta_x^i}{n} \\
\mathbb{Q} &= \sum_{i=1}^n \frac{\delta_y^i}{n} \\
H[\pi] &= \sum_{i=1}^n \sum_{j=1}^n \pi_{ij} \log \pi_{ij}
\end{aligned}$$

For **continuous** distributions:

$$\min_{\pi \in U[\mathbb{P}, \mathbb{Q}]} \int c[x, y] d\pi[x, y] + \varepsilon \int d\pi \log \frac{\pi[x, y]}{\mathbb{P}[x] \mathbb{Q}[y]} \quad (1)$$

**Notes:**

- $U[\mathbb{P}, \mathbb{Q}]$  is the set of couplings of  $\mathbb{P}$  and  $\mathbb{Q}$
- The log element in the above equation is equal to

$$KL[\pi | \mathbb{P}\mathbb{Q}]$$

**Entropic Optimal Transport:**

$$\min_{\pi \in U[\mathbb{P}, \mathbb{Q}]} \int c[x, y] d\pi[x, y] + \varepsilon \int d\pi \log \frac{\pi[x, y]}{\mathbb{P}[x] \mathbb{Q}[y]} \quad (2)$$

$$\begin{aligned} \text{where } \log \frac{\pi[x, y]}{\mathbb{P}[x] \mathbb{Q}[y]} &= KL[\pi | \mathbb{P}\mathbb{Q}] \\ \pi &\sim [x, y] \end{aligned}$$

- In information theory, this is called **mutual information**

**Claim:** the above equation is equivalent to

$$\min_{\pi \in U[\mathbb{P}, \mathbb{Q}]} KL[\pi | R] \quad (3)$$

$$\text{where } dR[x, y] = e^{-\frac{c[x, y]}{\varepsilon}} d\mathbb{P}[x] d\mathbb{Q}[y]$$

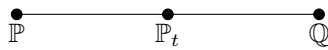
- **Note:** if we have a state space,  $X$  then  $R$  is a measure of  $XX$

**Proof of equivalence:**

$$\begin{aligned} &\int c[x, y] d\pi[x, y] + \varepsilon \int d\pi \log \frac{\pi[x, y]}{\mathbb{P}[x] \mathbb{Q}[y]} \\ &= \varepsilon \int [\log[e^{\frac{c}{\varepsilon}}] + \log \frac{\pi[x, y]}{\mathbb{P}[x] \mathbb{Q}[y]}] d\pi \\ &= \varepsilon \frac{\pi}{\mathbb{P}\mathbb{Q} e^{-\frac{c}{\varepsilon}}} d\pi \\ &= \varepsilon KL[\pi | R] \end{aligned}$$

**Dynamical Formulation of Optimal Transport**

Without entropy: we saw that OT coupling gives a shortest path in  $W_z$  connecting  $\mathbb{P}$  to  $\mathbb{Q}$ . We can interpolate along the constant speed geodesic  $\mathbb{P}_t$  for  $t \in (0, 1)$ .



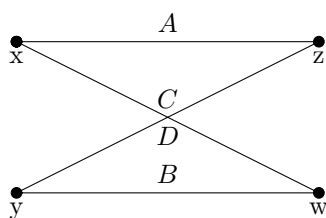
To get  $\mathbb{P}_t$ , take  $\pi^*$  and apply the map

$$[x, y] \mapsto [1 - t]x + ty$$

$$[x, y] \sim \pi$$

$$\text{law}[[1 - t]x + ty] = \mathbb{P}_t$$

With entropy:



1.  $\varepsilon$  low
  - 99% A and 1% C
2.  $\varepsilon$  medium
  - 90% A and 10% C
3.  $\varepsilon$  high
  - 50% A and 50% C