Tuesday, September 24 Prafi\_ "Let g= (a,m)\_ If ax = 1 lmom the >x = 1 (mod g) Shee glms but · Today: Group Work #2 Since  $g/\partial$ ,  $D=D_X \equiv d_{X} \equiv 1 \pmod{g}$ , (solutions posted ofter class) · Homework #2 posted today, due next Tuesday Forcing g=1. • Two quick proofs of enistence when (2,m)=1: Definition: Given mein and ac Z, - By Enler's theorem, we call XE I a (multiplicative) inverse of a modulo m if ax = 1 (mod m)-Example: 6.73 = 438 = 1 (mod 437) - By Bezout's theorem, we can unke 50 73 is on inverse of 6 (mo) 437) >x+ny=1 for some xy; then Theorem . If (2, m)>1 then a  $\partial x = \partial x + Oy \equiv \partial x + my = 1 (mo) m)$ - Uniquenest: if  $\partial x \equiv 1 (mo) m \partial x^2$ wis no inverse Lmod m?. • If (D,m)=1 then 2 has an inverse (modm), which is chique is a residue class (mod m). We denote this inverse by a (mod m).

Définition: If kell, we define Some proof - (gm)=1. // a long m) to be la' J (mom), when Corollory: If at 5 c (mod m) and (5m)=1, then (0,m)=1. Com)=1. Exercise. check everything Lemmis: Given MEIN and XEZ, 5211 mostes\_ consider these statements: Exercise: Suppose (2,m)=1 and  $(L_{1}, H_{m})=1$ . Let  $l = k^{2}$  (mod  $(L_{m})$ ).  $20 \times 21 \pmod{m^{2}}$  $(2) \chi^{-1} \equiv \chi \pmod{m};$ The (ak) = a (and m). (3)  $X \equiv \pm 1$  (mod m). - Retated to RSA cryptography · For any mEIN, (1) (2) and (3) => (1). Miscoloneous Lemmi - If ab = c (modm) · If m is probe, the (1) = (3) and (C,m)=1, then (D,m) = (b,m)=1. (so su three are equivalent) Prof: Wale c-ab=mn. The c= abturn is a multiple of Cann); so Cann divides both m ond c, Example:  $q^2 = 81 \in 1 \pmod{20}$ , so  $q^{-1} = 9 \pmod{20}$ ; but 9== ±1 (mor) 207. hence tom? [ Cym]=1. => Loum]=1.