

Tuesday, January 28

Group Work #3 in class on Thursday

RECALL:

$$\boxed{\eta(s) = \frac{1}{s-1} + 1 - s \int_1^{\infty} \{u\} u^{-s-1} du.}$$

Corollary 1.14: For $0 < s < 1$,

$$\frac{1}{s-1} < \eta(s) < \frac{s}{s-1} < 0.$$

Proof: Since $0 \leq \{u\} < 1$, we have

$$0 \leq s \int_1^{\infty} \{u\} u^{-s-1} du \leq s \int_1^{\infty} 1 \cdot u^{-s-1} du \\ = -u^{-s} \Big|_1^{\infty} = 1.$$

So $1-s \in [0, 1]$, so

$$\eta(s) = \frac{1}{s-1} + (1-s \int_1^{\infty} \{u\} u^{-s-1} du) \in \left[\frac{1}{s-1}, \frac{1}{s-1} + 1 \right],$$

• let's get an asymptotic formula for

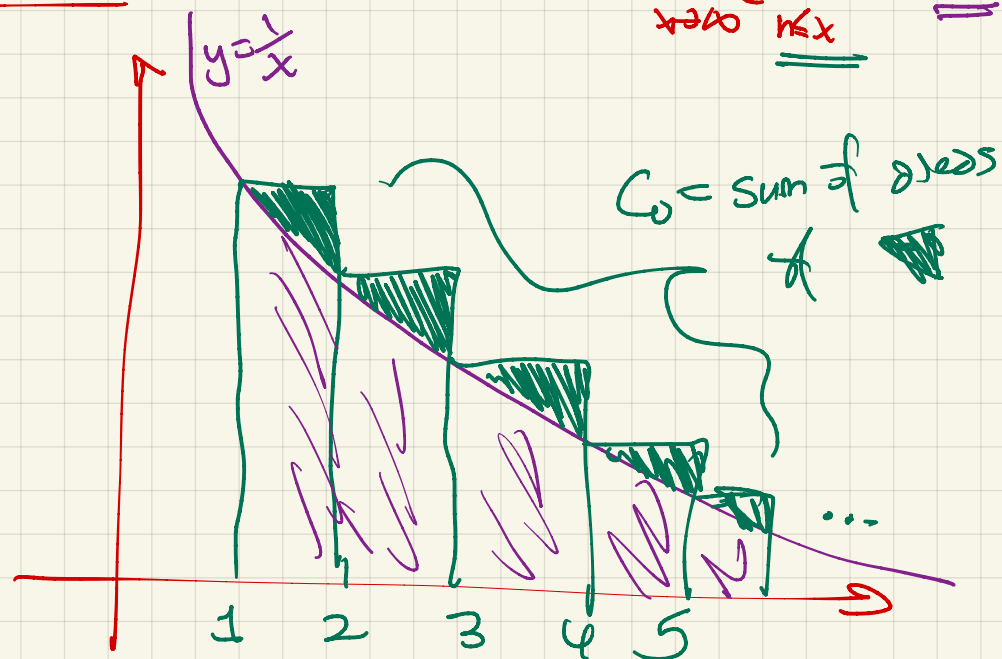
$$\begin{aligned} \sum_{n \leq x} \frac{1}{n} &= \int_1^x \frac{1}{u} d|u| \\ &= \int_1^x \frac{1}{u} du - \int_1^x \frac{1}{u} d\{u\} \\ &= \log u \Big|_1^x - \left(\frac{1}{u} \{u\} \Big|_1^x \right) + \int_1^x \{u\} d\left(\frac{1}{u}\right) \\ &= \log x - 0 - \left(\frac{\{x\}}{x} - \frac{1}{1} \right) - \int_1^x \frac{\{u\}}{u^2} du. \end{aligned}$$

$$\begin{aligned} \sum_{n \leq x} \frac{1}{n} &= \log x + O\left(\frac{1}{x}\right) + 1 - \int_1^{\infty} \frac{\{u\}}{u^2} du \\ &\quad + \int_x^{\infty} \frac{\{u\}}{u^2} du \\ &= \log x + C_0 + O\left(\frac{1}{x}\right). \end{aligned}$$

↙ Euler's constant

Note: $O\left(\frac{1}{x}\right)$ is the best error we can hope for,
since $\sum_{n \leq x} \frac{1}{n}$ has jumps of that size.

Note 2: We see that $C_0 = \lim_{x \rightarrow \infty} \left(\sum_{n \leq x} \frac{1}{n} - \log x \right)$



Note: picture shows $\frac{1}{2} < C_0 < 1$
(in fact $C_0 \approx 0.577$).

Conjecture C_0 is irrational

Chapter 2: Elementary estimates for arithmetic functions.

Motivating example: $0 < \frac{\phi(n)}{n} \leq 1$;
What's the average value (expectation)?

• Limit to $n \leq x$: $\frac{1}{x} \sum_{n \leq x} \frac{\phi(n)}{n}$.

• If $\lim_{x \rightarrow \infty} \left(\frac{1}{x} \sum_{n \leq x} \frac{\phi(n)}{n} \right)$ exists, we call it the mean value of $\frac{\phi(n)}{n}$.

Recall: $\phi \# 1 = n$; & $\phi = n * \mu$:

$$\phi(n) = \sum_{d|n} \mu(d) \frac{n}{d} \Leftrightarrow \frac{\phi(n)}{n} = \sum_{d|n} \frac{\mu(d)}{d}.$$

Then:

...

$$\begin{aligned}
\frac{1}{x} \sum_{n \leq x} \frac{\phi(n)}{n} &= \frac{1}{x} \sum_{n \leq x} \sum_{d|n} \frac{\mu(d)}{d} \\
&= \frac{1}{x} \sum_{d \leq x} \frac{\mu(d)}{d} \sum_{\substack{n \leq x \\ d|n}} 1 \\
&= \frac{1}{x} \sum_{d \leq x} \frac{\mu(d)}{d} \left\lfloor \frac{x}{d} \right\rfloor = \frac{1}{x} \sum_{d \leq x} \frac{\mu(d)}{d} \left(\frac{x}{d} + O(1) \right) \\
&= \sum_{d \leq x} \frac{\mu(d)}{d^2} + O\left(\frac{1}{x} \sum_{d \leq x} \frac{|\mu(d)|}{d} \right) \\
&= \sum_{d=1}^{\infty} \frac{\mu(d)}{d^2} + O\left(\sum_{d > x} \frac{|\mu(d)|}{d^2} + \frac{1}{x} \sum_{d \leq x} \frac{|\mu(d)|}{d} \right) \\
&= \frac{1}{6} + O\left(\sum_{d > x} \frac{1}{d^2} + \frac{1}{x} \sum_{d \leq x} \frac{1}{d} \right) \\
&= \frac{1}{\pi^2/6} + O\left(\frac{1}{x} + \frac{\log x}{x} \right) \\
&= \frac{6}{\pi^2} + O\left(\frac{\log x}{x} \right).
\end{aligned}$$

Hence $\lim_{x \rightarrow \infty} \frac{1}{x} \sum_{n \leq x} \frac{\phi(n)}{n} = \frac{6}{\pi^2}$

is the mean value of $\phi(n)/n$. ✓

General idea: If $f = g * 1$, then

$$\begin{aligned}
\sum_{n \leq x} f(n) &= \sum_{n \leq x} \sum_{d|n} g(d) \\
&= \sum_{d \leq x} g(d) \sum_{\substack{n \leq x \\ d|n}} 1 = \sum_{d \leq x} g(d) \left(\frac{x}{d} + O(1) \right)
\end{aligned}$$

Example: let's estimate $Q(x)$, the number of squarefree integers up to x :

$$Q(x) = \sum_{n \leq x} \begin{cases} 1, & \text{if } n \text{ is squarefree} \\ 0, & \text{otherwise} \end{cases}$$

$$= \sum_{n \leq x} \mu(n)^2.$$

Exercise: show that $\mu^2 = 1 * g$ where

" $g(d) = \mu(2\sqrt{d})$ ", that is,

$$g(d) = \begin{cases} \mu(d), & \text{if } d=c^2, \\ 0, & \text{if } d \text{ is not a square.} \end{cases}$$

So

$$(Q(x) = \sum_{n \leq x} \mu^2(n) = \sum_{d \leq x} g(d) \left(\frac{x}{d} + O(1) \right)$$

$$\stackrel{d=c^2}{=} \sum_{c \leq \sqrt{x}} \mu(c) \left(\frac{x}{c^2} + O(1) \right)$$

$$= x \sum_{c \leq \sqrt{x}} \frac{\mu(c)}{c^2} + O\left(\sum_{c \leq \sqrt{x}} |\mu(c)|\right)$$

$$= x \left(\sum_{c=1}^{\infty} \frac{\mu(c)}{c^2} + O\left(\sum_{c > \sqrt{x}} \frac{|\mu(c)|}{c^2}\right) \right) + O(\sqrt{x})$$

$$= x \left(\frac{6}{\pi^2} + O\left(\frac{1}{\sqrt{x}}\right) \right) + O(\sqrt{x})$$

$$= \frac{6}{\pi^2} x + O(\sqrt{x}).$$

"The probability that a randomly chosen integer is squarefree is $6/\pi^2$."

Terminology: Given $A \subseteq \mathbb{N}$, define its (natural) density to be

$$\delta(A) = \lim_{x \rightarrow \infty} \frac{1}{x} \# \{a \in A : a \leq x\},$$

if the limit exists.

Example: $A = \{a \in \mathbb{N} : a \equiv b \pmod{c}\}$, then

$$\# \{a \in A : a \leq x\} = \frac{x}{c} + O(1); \text{ so}$$

$$\delta(A) = \frac{1}{c}.$$

We showed $\delta(\{\text{squarefree numbers}\}) = \frac{6}{\pi^2}.$

$$= \frac{1}{\zeta(2)} = \prod_p \left(1 - \frac{1}{p^2}\right).$$

$$\delta(\{\text{squarefree numbers}\}) = \frac{6}{\pi^2}.$$

$$= \frac{1}{\delta(2)} = \prod_p \left(1 - \frac{1}{p^2}\right).$$

Note: it's ^{pretty} easy to show that for any

fixed y ,

$$\delta\left(\left\{n \in \mathbb{N}: p \leq y \Rightarrow p^2 \nmid n\right\}\right) = \prod_{p \leq y} \left(1 - \frac{1}{p^2}\right)$$

In fact,

$$\begin{aligned} \#\{n \leq x: p \leq y \Rightarrow p^2 \nmid n\} \\ = x \prod_{p \leq y} \left(1 - \frac{1}{p^2}\right) + O_y(1) \end{aligned}$$

But we can't take $\lim_{y \rightarrow \infty}$ because of the O_y .

Leveling up the method: Let $A(x) = \sum_{n \leq x} a_n$

and $B(x) = \sum_{n \leq x} b_n$. Define $c = a * b$

and write $C(x) = \sum_{n \leq x} c_n$. Then

$$C(x) = \sum_{n \leq x} \sum_{d|n} a_d b_{n/d}$$

$$= \sum_{d \leq x} a_d \sum_{\substack{n \leq x \\ d|n}} b_{n/d} = \sum_{d \leq x} a_d \sum_{\substack{n=md \\ m \leq \frac{x}{d}}} b_m$$

$$= \sum_{d \leq x} a_d B\left(\frac{x}{d}\right).$$

Previous example: $b_n \equiv 1$, so $B(x) = 2x + O(1)$

Example: let $c_n = d(n)$, so that $c = 1 * 1$.

$$\text{Example } \sum_{n \leq x} d(n).$$

$$C(x) = \sum_{d \leq x} a_d B\left(\frac{x}{d}\right).$$

Here, $a_d \equiv 1 \Rightarrow B\left(\frac{x}{d}\right) = \left\lfloor \frac{x}{d} \right\rfloor$, so

$$\begin{aligned} \sum_{n \leq x} d(n) &= \sum_{d \leq x} 1 \cdot \left(\frac{x}{d} + O(1)\right) \\ &= x \sum_{d \leq x} \frac{1}{d} + O\left(\sum_{d \leq x} 1\right) \\ &= x \left(\log x + C_0 + O\left(\frac{1}{x}\right)\right) + O(x) \\ &= x \log x + O(x). \end{aligned}$$

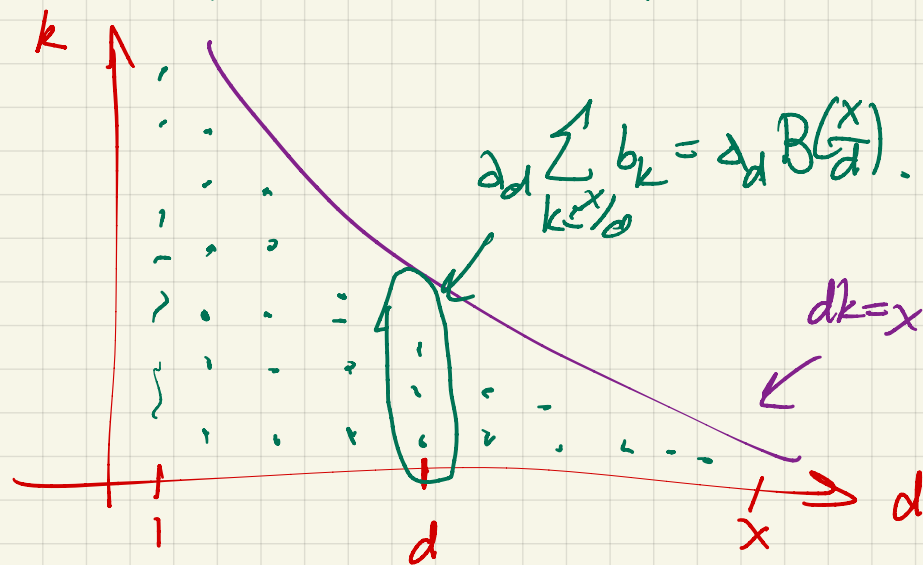
Exercise: For $x \geq 1$, $\sum_{n \leq x} \log n = x \log x + O(x)$.

Thus $\sum_{n \leq x} d(n) \sim \sum_{n \leq x} \log n$

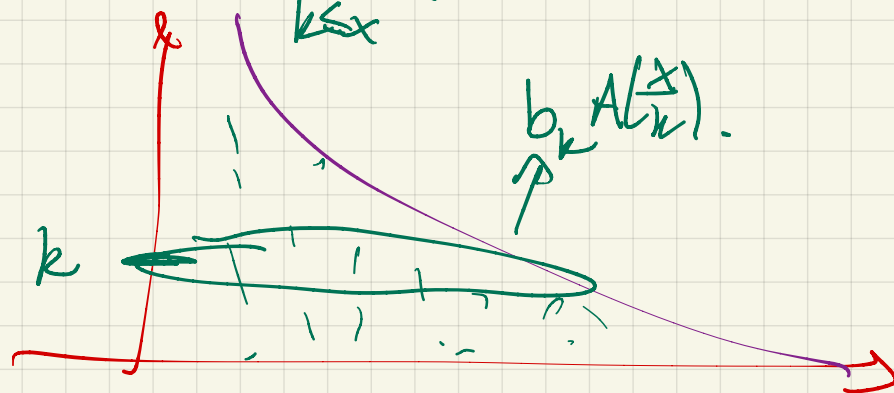
"The average order of $d(n)$ is $\log n$."

Why is the error so big?

$$\sum_{n \leq x} \sum_{d|n} a_d b_k = \sum_{d \leq x} a_d B\left(\frac{x}{d}\right).$$



Also, $\sum_{k \leq x} b_k A\left(\frac{x}{k}\right)$.



Clear ideas:-

