

Wednesday, March 15

Warm-up calculation: Recall

$$\psi(x, \chi) = \sum_{n \leq x} \Delta(n) \chi(n).$$

We can work out

$$\frac{1}{\phi(q)} \sum_{x \pmod q} |\psi(x, \chi)|^2 = \frac{1}{\phi(q)} \sum_{x \pmod q} \left( \sum_{n \leq x} \Delta(n) \chi(n) \right) \overline{\left( \sum_{m \leq x} \Delta(m) \chi(m) \right)}$$

$$= \frac{1}{\phi(q)} \sum_{\substack{n \leq x \\ m \leq x}} \Delta(n) \Delta(m) \sum_{x \pmod q} \chi(n) \overline{\chi(m)}$$

$$= \sum_{\substack{n \leq x \\ m \leq x \\ (nm, q) = 1 \\ n \equiv m \pmod q}} \Delta(n) \Delta(m) =$$

$$= \sum_{\substack{a \pmod q \\ (a, q) = 1}} \sum_{\substack{n \leq x \\ m \leq x \\ n \equiv m \pmod q}} \Delta(n) \Delta(m)$$

$$= \sum_{\substack{a \pmod q \\ (a, q) = 1}} \left( \sum_{\substack{n \leq x \\ n \equiv a \pmod q}} \Delta(n) \right)^2$$

$$= \sum_{\substack{a \pmod q \\ (a, q) = 1}} \chi(x; q, a)^2.$$

Exercise: Show that

$$V(x; q) = \frac{1}{\phi(q)} \sum_{\substack{x \pmod q \\ x \neq x_0}} |\psi(x, \chi)|^2$$

$$= \sum_{\substack{a \pmod q \\ (a, q) = 1}} \left| \psi(x; q, a) - \frac{1}{\phi(q)} \sum_{n \leq x} \Delta(n) \right|^2$$

and also

$$G(x; q) = \sum_{(x, q) = 1} \left| \theta(x; q, \alpha) - \frac{x}{\phi(q)} \right|^2$$
$$= \frac{1}{\phi(q)} \sum_{x \pmod{q}} \left| \theta(x; q) - \begin{cases} x, & \text{if } x \equiv \alpha, \\ 0, & \text{otherwise} \end{cases} \right|^2$$

Huxley (1970s) conjectured that as soon as  $q$  goes to infinity with  $x$ , we have

$$\frac{G(x; q)}{V(x; q)} \ll x \log q.$$

Note: we don't expect this upper bound when  $q$  is fixed.

• Littlewood (1910s):

$$\pi(x; 4, 3) - \pi(x; 4, 1) = \mathcal{O}_{\epsilon} \left( \frac{\sqrt{x}}{\log x} (\log \log \log x) \right)$$
$$\pi(x; 3, 2) - \pi(x; 3, 1)$$

which implies

$$G(x; 4)$$

$$V(x; 4) = \mathcal{O} \left( x (\log \log \log x)^2 \right)$$

$$G(x; 3)$$

$$V(x; 3)$$

which is not  $\ll x$ .

- Davidoff (1980s?) showed for fixed  $q$ ,

$$\frac{G(x; q)}{V(x; q)} = \mathcal{O} \left( x (\log \log \log x)^2 \right).$$

not  $\ll x$ .

The average version,

$$(*) \quad \frac{1}{Q} \sum_{q \leq Q} V(x; q) \ll x \log Q,$$

is the Barban-Davenport-Halberstam Thm,  
valid for  $Q = x / (\log x)^A$ .

$$V(x; q) = \frac{1}{\phi(q)} \sum_{\substack{x \pmod{q} \\ x \neq x_0}} |\psi(x; q)|^2$$

- For individual  $q$ : if we assume GRH and a strong version of the prime pairs conjecture (Hardy-Littlewood), Friedlander-Goldston showed  $V(x; q) \ll x \log q$  for  $q \geq x^{\frac{1}{2} + \varepsilon}$ .

- For  $q < x^{\frac{1}{2}}$ , nothing known (unless conditionally).

Fiorilli (2015), "The distribution of the variance of primes in arithmetic progressions". conjectured:

$$V(x; q) \ll x \log q \text{ when } q > (\log \log x)^{1+\varepsilon}$$

but not when  $q < (\log \log x)^{1-\varepsilon}$ .

Heuristic/motivation:

There's some random variable  $H_q$  that models  $V(x; q)$  (or  $\phi(q)$ ).

- recall that

$$\psi(x; q) = - \sum_{\substack{p \\ (p, q) = 1 \\ p \leq x}} \frac{x}{p} + \text{small}$$

$$\mathbb{E}(H_q) \sim \phi(q) \log q$$

$$\sigma^2(H_q) \sim 2\phi(q)(\log q)^2$$

$H_q$  "roughly normal variable"

models  $\phi(q) \frac{V(x; q)}{x}$

Fiorilli proved, about  $H_q$ , that

$$\begin{aligned} \frac{1}{4}e^{-c_1\varepsilon^2\phi(q)} &\leq \\ \Pr(|H(q) - \phi(q)\log q| > \varepsilon\phi(q)\log q) & \\ &\leq 2e^{-c_2\varepsilon^2\phi(q)} \end{aligned}$$

(for some  $c_1, c_2 > 0$ ).

These are "tail estimates" for  $H_q$ .

⇒ heuristically give conjectures  
for how large  $x$  needs to be  
for large deviations to occur.

⇒ Fiorilli's cutoff  $\log \log x$

Fiorilli-M. (2023+) "Disproving  
Hooley's conjecture"

$$\text{Theorem: } \frac{G(x;q)}{V(x;q)} = \sum L\left(x \log q, \frac{\log \log x}{2}\right)^*$$

In particular, Hooley's conjecture  
fails to hold in the range

$q < \delta \log \log x$  if  $\delta$  is  
sufficiently small.

We proved \* holds for a positive  
proportion of  $q$ .

- GRH is false: \* holds for every multiple  $g$  of  $g_0$ , where  $\text{GRH}(g_0)$  is false
- GRH is true: extension of the (All) Hardy-Littlewood / Dirichlet method ( $q$ )