

I
Embedding of Algebras

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1. Abstract

We look at some general conditions for determining when the unit of an adjunction is monic when the generated monad is defined on a category of algebras. Among other possible applications these conditions will be shown to provide a common approach to some classic embedding theorems in algebra.

2. Introduction

Let $T = (T, \eta, \mu)$ be a monad defined on a category \mathbf{C} . We consider the following problem. Is η_C monic for a given object C of \mathbf{C} ?

There is an easy answer in the case of monads defined on \mathbf{Set} . This is given here but goes back at least to the thesis of Manes.

In this talk we require that T be defined on a category \mathbf{C} of Eilenberg-Moore algebras given by finitary operations and equations. It is further required that T be generated by a pair of adjoint functors fitting into a certain framework.

Necessary and sufficient conditions are given for η_C to be monic for given object C when T fits into the framework described. A key condition is given in terms of a graph assigned to object C .

Additional flexibility is attained by relating the key condition to several other graph conditions. This results in another theorem, which we apply in several cases.

3. Monads over \mathbf{Set} .

Let $T = (T, \eta, \mu)$ be a monad defined on \mathbf{Set} . The category \mathbf{Set}^T of Eilenberg-Moore algebras is said to be *nontrivial* if there is at least one algebra (A, α) whose underlying set A has more than one element.

Lemma. *If the category of Eilenberg-Moore algebras for a monad (T, η, μ) on \mathbf{Set} is nontrivial, then η_X is monic for all sets X .*

Proof. Let x and y be distinct elements of a set X with more than one element and suppose that (A, α) is an algebra with more than one element. Then there is a function $f : X \rightarrow A$ with $f(x)$ distinct from $f(y)$. By the universal mapping property there is a unique algebra morphism $g : (TX, \mu_X) \rightarrow (A, \alpha)$ such that the following diagram commutes

$$\begin{array}{ccc} X & \xrightarrow{\eta_X} & TX \\ & \searrow f & \downarrow g \\ & & A \end{array}$$

thus $\eta_X(x)$ is not equal to $\eta_X(y)$. □

Thus for example, in the case of semigroups, this shows that any set may be embedded in the free semigroup.

It is clearly rare for the category of Eilenberg-Moore algebras to be trivial since, if trivial, each TX is either

empty or consists of one element since it is the underlying object of the algebra (TX, μ_X) . Furthermore, TX is the codomain of η_X , hence cannot be empty unless X is.

Now, we remark that the unique function $T1 \rightarrow 1$ is a T -algebra, where 1 is the one point set, and the empty set has a T -structure if and only if $T(\emptyset) = \emptyset$. Accordingly, up to isomorphism, there are only two monads (T_1, η_1, μ_1) and (T_2, η_2, μ_2) yielding just trivial algebras. These are given by $T_1X = 1$ for each set X and $T_2(\emptyset) = \emptyset$ otherwise $T_2(X) = 1$

4. Reduction and Graph Components.

Let \mathbf{G} be a small subcategory of a category \mathbf{C} and let $\mathcal{P}(\mathbf{G})$ be the *power category* of \mathbf{G} . The objects of $\mathcal{P}(\mathbf{G})$ are the subclasses of objects of \mathbf{G} and the morphisms are the inclusions.

The *reduction functor* $\mathcal{R}_{\mathbf{G}} : \mathbf{C}^{\text{op}} \rightarrow \mathcal{P}(\mathbf{G})$ is defined by

$$\mathcal{R}_{\mathbf{G}}X = \{A \mid A \text{ is in } \mathbf{G} \text{ and } X \rightarrow A \text{ exists in } \mathbf{C}\}$$

with the obvious definition on morphisms.

An object A is *reduced* in \mathbf{C} if the only \mathbf{C} -morphism with domain A is the identity. The subcategory \mathbf{G} is *reduced in* \mathbf{C} if each of its objects is reduced in \mathbf{C} . An object X of \mathbf{C} is *\mathbf{G} -reducible* if $\mathcal{R}_{\mathbf{G}}X$ is nonempty.

The *component class* $[X]$ of an object X of a category \mathbf{C} (or a graph \mathbf{C}) is the class of all objects Y which can be connected to X by a finite sequence of morphisms (e.g. $X \rightarrow X_1 \longleftarrow X_2 \rightarrow Y$). We let $\text{Comp}\mathbf{C}$ denote the collection of component classes.

Strong Embedding Principle for \mathbf{G} If $[X] = [Y]$ in $\text{Comp}\mathbf{C}$, then $\mathcal{R}_{\mathbf{G}}X = \mathcal{R}_{\mathbf{G}}Y$. Furthermore there is at most one morphism $X \rightarrow A$ for each pair (X, A) consisting of an object X of \mathbf{C} and an object A of \mathbf{G} .

5. A Framework for Embeddings.

Let $\mathbf{Alg}(\Omega, E)$ denote the category of algebras defined by a set of operators Ω and identities E . Suppose

$$\mathbf{A} \xrightarrow{U} \mathbf{B} \xrightarrow{V} \mathbf{D} \quad (1)$$

is a diagram such that

- (a) \mathbf{A} , \mathbf{B} and \mathbf{D} are the categories of (Ω, E) , (Ω', E') and (Ω'', E'') algebras, respectively, with $\Omega'' \subseteq \Omega'$ and $E'' \subseteq E'$, and
- (b) V is the forgetful functor on operators $\Omega' - \Omega''$ and identities $E' - E''$ and U is a functor commuting with the underlying set functors on \mathbf{A} and \mathbf{B} . Note that U is *not* necessarily a functor forgetting part of Ω and E .

We next describe a functor $C_V : \mathbf{B} \rightarrow \mathbf{Grph}$ associated to each pair consisting of a diagram (1) of algebras and an adjunction $(L, VU, \varphi') : \mathbf{D} \rightarrow \mathbf{A}$, where \mathbf{Grph} is the category of directed graphs.

Given an object G of \mathbf{B} let the objects of the graph $C_V(G)$ be the elements of the underlying set $|LVG|$ of LVG .

Recursive definition of the arrows of $C_V(G)$:

$$\omega_{ULVG}(|\eta'_{VG}|x_1, \dots, |\eta'_{VG}|x_n) \rightarrow |\eta'_{VG}|\omega_G(x_1, \dots, x_n)$$

is an arrow if ω is in the set $\Omega' - \Omega''$ of operators forgotten by V and (x_1, \dots, x_n) is an n-tuple of elements of $|G|$ for which $\omega_G(x_1, \dots, x_n)$ is defined.

If $d \rightarrow e$ is an arrow of $C_V(G)$, then so is

$$\rho_{LVG}(d_1, \dots, d, \dots, d_q) \rightarrow \rho_{LVG}(d_1, \dots, e, \dots, d_q)$$

for ρ an operator of arity q in Ω and $d_1, \dots, d_{i-1}, d_{i+1}, \dots, d_q$ arbitrary elements of $|LVG|$.

If $\beta : G \rightarrow G' \in \mathbf{B}$, then $C_V(\beta) : C_V(G) \rightarrow C_V(G')$ is the graph morphism which is just the function $|LV\beta| : |LVG| \rightarrow |LVG'|$ on objects and defined recursively on arrows in the obvious way.

In the following proposition note that if, in the diagram (1), \mathbf{D} is the category of sets, then an adjunction (L, VU, φ') is given by letting LX be the free (Ω, E) algebra on the set X .

Proposition. *Suppose*

$$\begin{array}{ccccc} \mathbf{A} & \xrightarrow{U} & \mathbf{B} & \xrightarrow{V} & \mathbf{D} \\ & & \searrow & \swarrow & \\ & & & L & \end{array}$$

is a diagram of algebras as in (1), with given adjunction $(L, VU, \varphi') : \mathbf{D} \rightarrow \mathbf{A}$. Then there is an adjunction $(F, U, \varphi) : \mathbf{B} \rightarrow \mathbf{A}$ with the following specific properties:

- (a) *The underlying set of FG is $\text{Comp}C_V(G)$.*
- (b) *If ρ is an operator of arity n in Ω , then ρ_{FG} is defined by*

$$\rho_{FG}([c_1], \dots, [c_n]) = [\rho_{LVG}(c_1, \dots, c_n)]$$

where c_1, \dots, c_n are members of the set $|LVG|$ of objects of the graph $C_V(G)$.

- (c) *The unit morphism $\eta_G : G \rightarrow UFG$ of (F, U, ϕ) has an underlying set map which is the composition $[-] \cdot |\eta'_{VG}|$, where $|\eta'_{VG}| : |G| \rightarrow |LVG| = \text{Obj}C_V(G)$ is the set map underlying the unit $\eta'_{VG} : VG \rightarrow VULVG$ of the adjunction (L, VU, ϕ') and $[-] : \text{Obj}C_V(G) \rightarrow \text{Comp}C_V(G)$ is the component function.*

Suppose the hypotheses of the proposition hold. Let S_V be a subgraph of $C_V(G)$ having the same objects $|LVG|$ and the same components as $C_V(G)$. Then the proposition remains valid under substitution of S_V for $C_V(G)$ throughout. This allows us to “picture” the adjoint using a possibly smaller set of arrows than those present in $C_V(G)$. Accordingly, we define a *V picture of the adjoint F to U at $G \in |\mathbf{B}|$* to be any quotient category $\mathbf{C}(= \mathbf{C}(S_V))$ of the free category generated by such a subgraph S_V of $C_V(G)$. This proposition is then

valid upon substitution of the underlying graph of a V picture \mathbf{C} for $C_V(G)$ throughout.

We now present the first embedding theorem.

Theorem. *Let*

$$\mathbf{A} \begin{array}{c} \xrightarrow{U} \\ \xleftarrow{F} \end{array} \mathbf{B} \xrightarrow{V} \mathbf{D} \begin{array}{c} \xleftarrow{L} \\ \xrightarrow{L} \end{array} \mathbf{A}$$

be given with adjunctions $(L, VU, \phi') : \mathbf{D} \rightarrow \mathbf{A}$ and $(F, U, \phi) : \mathbf{B} \rightarrow \mathbf{A}$ as described in the Proposition.

Given $G \in |\mathbf{B}|$ let $\mathbf{C}(S_V)$ be any V picture of the adjoint F to U at G .

Then the unit morphism $\eta_G : G \rightarrow UFG$ of the adjunction (F, U, ϕ) is monic if and only if the following hold

- (a) The discrete subcategory $\mathbf{G} = \eta'_{VG}(|G|)$ is reduced in $\mathbf{C}(S_V)$ for η'_{VG} the unit of (L, VU, ϕ') .*
- (b) If $[A] = [B]$ in $\text{Comp}\mathbf{C}(S_V)$ with $A, B \in |\mathbf{G}|$, then $\mathcal{R}_{\mathbf{G}}A = \mathcal{R}_{\mathbf{G}}B$, where $\mathcal{R}_{\mathbf{G}} : \mathbf{C}(S_V)^{\text{op}} \rightarrow \mathcal{P}(\mathbf{G})$ is the reduction functor.*
- (c) The unit morphism η'_{VG} is monic.*

6. Diamonds, Strong Embeddings and Connectedness.

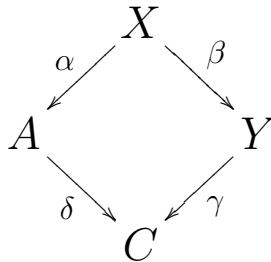
We make use of the following principles.

Diamond Principle for \mathbf{G} . Let \mathbf{G} be a subcategory of \mathbf{C} , then the objects $\alpha : X \rightarrow A$ of X/\mathbf{C} with A an object of \mathbf{G} are terminal in X/\mathbf{C} for each object X of \mathbf{C} .

Lemma A. *Let \mathbf{G} be a reduced subcategory of \mathbf{C} . Then the following statements are equivalent:*

- (a) *The Diamond Principle for \mathbf{G} .*
- (b) *Each pair $Y \leftarrow X \rightarrow Z$ of \mathbf{C} -morphisms with X a \mathbf{G} -reducible object can be completed to a commutative diamond in \mathbf{C} .*

Proof. Suppose that (b) holds and let $\alpha : X \rightarrow A$ be an object of X/\mathbf{C} with A in \mathbf{G} and $\beta : X \rightarrow Y$ be any other object. Then, by hypothesis there exists a commutative diagram



in \mathbf{C} . But A in \mathbf{G} implies that $\delta = 1$. Thus $\gamma : \beta \rightarrow \alpha$. If $\gamma' : \beta \rightarrow \alpha$ then, by hypothesis, $A \xleftarrow{\gamma} Y \xrightarrow{\gamma'} A$ can be completed to a commutative square since Y is G -reducible. Thus $\gamma = \gamma'$ since A is reduced and so α is terminal.

It is trivial to show that (a) implies (b). \square

For the next Lemma we recall the definition of the

Strong Embedding Principle for \mathbf{G} . If $[X] = [Y]$ in $\text{Comp}\mathbf{C}$, then $\mathcal{R}_{\mathbf{G}}X = \mathcal{R}_{\mathbf{G}}Y$. Furthermore there is at most one morphism $X \rightarrow A$ for each pair (X, A) consisting of an object X of \mathbf{C} and an object A of \mathbf{G} .

Lemma B. *The Strong Embedding Principle for \mathbf{G} implies the Diamond Principle for \mathbf{G} .*

Proof. Let $A \xleftarrow{\alpha} X \xrightarrow{\beta} Y$ be a diagram in \mathbf{C} with A in \mathbf{G} . Thus $[A] = [Y]$ in $\text{Comp}\mathbf{C}$ and $\mathcal{R}_{\mathbf{G}}X = \mathcal{R}_{\mathbf{G}}Y$, by hypothesis. Hence A is in $\mathcal{R}_{\mathbf{G}}Y$ and there exists $\gamma : Y \rightarrow A$. But then $\gamma\beta$ and α are morphisms $X \rightarrow A$ and $\alpha = \gamma\beta$ by hypothesis. Thus $\gamma : \beta \rightarrow \alpha$ in X/\mathbf{C} and γ is unique since as a \mathbf{C} morphism it is the only morphism $Y \rightarrow A$ by hypothesis. \square

The implication also goes in the other direction. For a proof see [6].

Given an object X of \mathbf{C} let $(X/\mathbf{C})_{\mathcal{P}}$ be the full sub-

category of the slice category X/\mathbf{C} obtained by omitting the object $1_X : X \rightarrow X$.

Principle of Connectedness for \mathbf{T} . The categories $(X/\mathbf{C})_{\mathcal{P}}$ are connected for each object X of \mathbf{T} , where \mathbf{T} is a subcategory of \mathbf{C} .

Lemma C. *If $\alpha : X \rightarrow A$ is a terminal object of X/\mathbf{C} with A reduced, then $(X/\mathbf{C})_{\mathcal{P}}$ is connected.*

Proof. If X is reduced, then X/\mathbf{C} contains only one object, namely 1_X , and $(X/\mathbf{C})_{\mathcal{P}}$ is empty, hence trivially connected. If X is not reduced and $\alpha : X \rightarrow A$ is terminal in X/\mathbf{C} with A reduced, then $\alpha \neq 1_X$ and α is terminal in $(X/\mathbf{C})_{\mathcal{P}}$. Thus $(X/\mathbf{C})_{\mathcal{P}}$ is connected. \square

Lemma D. *If the Diamond Principle holds for a reduced subcategory \mathbf{G} of \mathbf{C} , then the Principle of Connectedness holds for the full subcategory $\mathbf{T}_{\mathbf{G}}$ of \mathbf{C} consisting of all G -reducible objects of \mathbf{C} .*

Proof. Let X be in $\mathbf{T}_{\mathbf{G}}$. Then there is a morphism $X \rightarrow A$ in \mathbf{C} with A in \mathbf{G} . By the Diamond Principle for \mathbf{G} the morphism $X \rightarrow A$ is terminal in X/\mathbf{C} . By Lemma C then $(X/\mathbf{C})_{\mathcal{P}}$ is connected. \square

Definition. Let \mathbb{N} be the preorder of nonnegative integers with $n \rightarrow m$ iff $n \geq m$. A *rank functor* for a category \mathbf{C} is a functor $R : \mathbf{C} \rightarrow \mathbb{N}$ with $R\alpha \neq 1$ whenever $\alpha \neq 1$.

Theorem. *Let \mathbf{C} be a category with rank functor given and \mathbf{G} a subcategory which is reduced in \mathbf{C} . Then the following are equivalent.*

- (a) The Principle of Connectedness for the full subcategory $\mathbf{T}_{\mathbf{G}}$ of \mathbf{C} determined by all \mathbf{G} reducible objects of \mathbf{C} .*
- (b) The Diamond Principle for \mathbf{G} .*
- (c) The Strong Embedding Principle for \mathbf{G} .*

7. The Second Embedding Theorem.

In the presence of a rank functor we have seen in Theorem 6 that the three principles are equivalent. We apply this Theorem to the hypotheses of Theorem 5 to obtain the following result:

Theorem. *Let*

$$\mathbf{A} \begin{array}{c} \xrightarrow{U} \\ \xleftarrow{F} \end{array} \mathbf{B} \xrightarrow{V} \mathbf{D}$$

\xleftarrow{L}

be given with hypotheses as in Theorem 6.

Then the unit morphism $\eta_G : G \rightarrow UFG$ of (F, U, φ) is monic provided there exists a V picture \mathbf{C} of the adjoint F to U at G for which the following conditions hold:

- (a) \mathbf{C} has a rank functor.
- (b) The discrete subcategory $\mathbf{G} = \eta'_{VG}(|G|)$ is reduced in \mathbf{C} for η'_{VG} the unit of (L, VU, φ') .
- (c) The categories $(X/\mathbf{C})_{\mathcal{P}}$ are connected for each $X \in |\mathbf{C}|$ which is \mathbf{G} reducible.
- (d) The unit morphism η'_{VG} of (L, VU, φ') is monic.

8. Associative Embedding of Lie Algebras.

Let $U : \mathbf{A} \rightarrow \mathbf{L}$ be the usual algebraic functor from associative algebras over K to Lie algebras over K , for K a commutative ring. That is, for $A \in |\mathbf{A}|$, UA is the same as A except that a new multiplication $[ab] = a.b - b.a$ replaces the associative multiplication $a.b$ of $A \in |\mathbf{A}|$.

It is well known that an adjunction $(F, U, \varphi) : \mathbf{L} \rightarrow \mathbf{A}$ exists. The question of embeddability of a Lie algebra in its universal associative algebra (i.e. the question as to whether the unit morphisms η_G of (F, U, φ) are monomorphisms) has been investigated by various authors (cf. Birkhoff[3] and Serre[11]). Not all Lie algebras can be so embedded (cf. Higgins[4]).

We demonstrate how such a question can be put in the context of the previous sections. Let $V : \mathbf{L} \rightarrow Mod_K$ be the functor forgetting the Lie multiplication. We then have the diagram

$$\begin{array}{ccccc} \mathbf{A} & \xrightarrow{U} & \mathbf{L} & \xrightarrow{V} & Mod_K \\ & & & \searrow & \\ & & & & L & \swarrow \\ & & & & & \end{array}$$

where the conditions (a) and (b) of diagram (1) hold. The existence of the adjunction $(L, VU, \varphi') : Mod_K \rightarrow \mathbf{A}$ is

assured on theoretical grounds, but can also be described explicitly as follows.

Given $G \in |\mathbf{L}|$ it is known that LVG is the tensor algebra of VG . Thus

$$LVG = \bigoplus_{n \geq 0} (\bigotimes_{i=1}^n (VG)).$$

Furthermore $\eta'_{VG} : VG \rightarrow K \oplus VG \oplus (VG \otimes VG) \oplus \dots$ is monic. In this section let \mathbf{G} be the discrete subgraph $\eta'_{VG}(|G|)$ of $C_V(G)$. Thus \mathbf{G} is a discrete subcategory of any V picture of F at G . Applying Theorem 5 the following Lie algebra embedding result holds.

Theorem A. *Let G be a Lie algebra and \mathbf{C} any V picture of F at G . Then a Lie algebra G can be embedded in its universal associative algebra FG if and only if the following hold:*

- (a) $[A] = [B]$ in $\text{Comp}\mathbf{C}$ implies that $\mathcal{R}_{\mathbf{G}}A = \mathcal{R}_{\mathbf{G}}B$ for all $A, B \in |\mathbf{G}|$ where $\mathcal{R}_{\mathbf{G}} : \mathbf{C}^{\text{op}} \rightarrow \mathcal{P}(\mathbf{G})$ is the reduction functor.
- (b) \mathbf{G} is a reduced subcategory of \mathbf{C} .

Similarly, by applying Theorem 7 and noting that η'_{VG} is monic we have the following sufficient conditions:

Theorem B. *A Lie algebra G can be embedded in its universal associative algebra FG if there exists any*

V picture \mathbf{C} of the adjoint F to U at G with the following properties.

- (a) *The categories $(X/\mathbf{C})_{\mathcal{P}}$ are connected for each $X \in |\mathbf{C}|$ which is \mathbf{G} reducible.*
- (b) *\mathbf{C} has a rank functor.*
- (c) *The discrete subcategory $\mathbf{G} = \eta'_{VG}(|G|)$ is reduced in \mathbf{C} for η'_{VG} the unit of (L, VU, ϕ') .*

Theorem (Birkhoff-Witt). *A Lie algebra G whose underlying module VG is free can be embedded in its universal associative algebra FG .*

Proof. The conditions of the previous theorem are to be verified for the following V picture of the adjoint F to U at G . Let \mathbf{C} be the preorder which is a quotient of the free category on the following subgraph S_V of $C_V(G)$. The objects of S_V are the elements of the free K module LVG on all finite strings $x_{i_1} \cdots x_{i_n}$ of elements from a basis $(x_i)_{i \in I}$ of the free K module VG . Given a well ordering of I we let the arrows of S_V be those of the form

$$k_i x_{i_1} \cdots x_{i_n} + \alpha \rightarrow$$

$k_i x_{i_1} \cdots x_{i_{j+1}} x_{i_j} \cdots x_{i_n} + k_i x_{i_1} \cdots [x_{i_j}, x_{i_{j+1}}] \cdots x_{i_n} + \alpha$
for $i_{j+1} < i_j$, $k_i \in K$, and α any element of LVG (not involving $x_{i_1} \cdots x_{i_n}$).

To show that the categories $(X/\mathbf{C})_{\mathcal{P}}$ are connected for each $X \in |\mathbf{C}|$ which is \mathbf{G} reducible, it turns out that the key idea is to show that for $c < b < a$ in I the objects

$$\beta : x_a x_b x_c \rightarrow x_b x_a x_c + [x_a, x_b] x_c$$

and

$$\gamma : x_a x_b x_c \rightarrow x_a x_c x_b + x_a [x_b, x_c]$$

can be connected in $((x_a x_b x_c)/\mathbf{C})_{\mathcal{P}}$. This is done by further reduction of the ranges of β and γ and use of the Jacobi identity and the identity $[x, y] = -[y, x]$.

The rank functor for \mathbf{C} is given as follows. Given $X = kx_{a_1} \cdots x_{a_n}$ let $R(X) = (R_n(X))$ be a sequence of nonnegative integers defined by $R_n(X) = \sum_{i=1}^n p_{a_i}$ where p_{a_i} is the number of x_{a_j} to the right of x_{a_i} with $a_j < a_i$ and $R_s(X) = 0$ for $s \neq n$. We extend by linearity to all elements of $LVG = |\mathbf{C}|$. If $X \rightarrow Y$ is an arrow, then $R(Y) < R(X)$ where the latter inequality means that $R_n(Y) < R_n(X)$ for n the largest integer with $R_n(Y) \neq R_n(X)$. Thus R extends to a rank functor.

Finally, we verify condition (c) by observing that any element of $|G|$ may be expressed in the form $\sum_{i \in I} k_i x_i$ in terms of the basis $(x_i)_{i \in I}$ of \mathbf{G} , which is regarded as a subset \mathbf{G} of LVG via the embedding η'_{VG} . From the preceding description of arrows of S_V there is no arrow

with domain an element of $\mathbf{G} = \eta'_{VG}(|G|)$. Thus \mathbf{G} is reduced in \mathbf{C} . \square

9. Sets with a partially defined binary operation and a result of Schreier.

We show how the following classical theorem follows from sections 5 and 7.

Theorem (Schreier). *If S is a common subgroup of the groups X and Y and if*

$$\begin{array}{ccc} S & \xrightarrow{\subseteq} & X \\ \subseteq \downarrow & & \alpha \downarrow \\ Y & \xrightarrow{\beta} & P \end{array} \quad (2)$$

is the pushout in the category of groups, then α and β are monomorphisms. The group P is referred to as the free product of X and Y with amalgamated subgroup S .

Let \mathbf{B} be the category of sets with a single partially defined binary operation. The diagram $X \longleftarrow S \longrightarrow Y$ of groups can be regarded as a diagram in \mathbf{B} and can be completed to a diagram

$$\begin{array}{ccc} S & \xrightarrow{\subseteq} & X \\ \subseteq \downarrow & & \gamma \downarrow \\ Y & \xrightarrow{\delta} & G \end{array} \quad (3)$$

commuting in \mathbf{B} where G is the disjoint union of X and Y with common subset S identified and $a.b$ defined if both

$a, b \in X$ or if both $a, b \in Y$, otherwise $a.b$ is undefined. The morphisms γ and δ are the obvious monomorphisms. The next Lemma and Proposition describe how this approach yields the Schreier Theorem.

Lemma. *The Schreier Theorem holds if in \mathfrak{B} the pushout codomain G in \mathbf{B} is embeddable in a group.*

Proof. Let $\iota : G \rightarrow P'$ be a monomorphism in \mathbf{B} with P' a group. Then

$$\begin{array}{ccc} S & \xrightarrow{\subseteq} & X \\ \subseteq \downarrow & & \downarrow \iota\gamma \\ Y & \xrightarrow{\iota\delta} & P' \end{array}$$

commutes in groups. Thus for some group homomorphism ϕ we have $\iota\gamma = \phi\alpha$ and $\iota\delta = \phi\beta$ since (2) is a pushout in groups. Thus α, β are monic since ι, γ and δ are. \square

Proposition. *Let G be as in the lemma. Then G is embeddable in a group.*

Proof. We embed G in a particular semigroup which turns out to be a group. Begin with the diagram

$$\begin{array}{ccccc} \mathbf{A} & \xrightarrow{U} & \mathbf{B} & \xrightarrow{V} & \mathit{Sets} \\ & & & \searrow & \\ & & & L & \swarrow & \\ & & & & & \end{array}$$

where \mathbf{A} is the category of semigroups(not necessarily with 1), U forgetful, V forgetful and (L, VU, φ') an adjunction. We then have a preorder \mathbf{C}_G which is a quotient of the free category \mathbf{F}_G generated by $C_V(G)$. Proposition 5 (which also holds for the category \mathbf{B} of partial algebras, see[6]) shows that an adjunction $\langle F, U, \varphi \rangle: \mathbf{B} \rightarrow \mathbf{A}$ exists and describes it. It is sufficient to show that the unit $\eta_G : G \rightarrow UFG$ of the adjunction is a monomorphism. By Theorem 7 it is sufficient to verify conditions (a) through (d). These conditions are trivial except for (c), which requires that the categories $(X/\mathbf{C})_{\mathcal{P}}$ be connected for each \mathbf{G} -reducible object X of \mathbf{C} . The objects of \mathbf{C}_G are elements of the free semigroup LVG on VG . An object X may be written as a string (a_1, \dots, a_n) of length $n \geq 1$ where $a_i \in VG$ for $i = 1, \dots, n$. It is sufficient to show that C_V arrows

$$\begin{array}{ccc}
 & (a_1, \dots, a_n) & \\
 \alpha \swarrow & & \searrow \beta \\
 (\dots, a_i a_{i+1}, \dots) & & (\dots, a_j a_{j+1}, \dots)
 \end{array}$$

regarded as $(X/\mathbf{C}_G)_{\mathcal{P}}$ objects can be connected by a finite sequence of morphisms in the same category. This requires a detailed argument when $i = j - 1$ or $i = j + 1$, otherwise it is trivial (cf. Baer[1], MacDonald[7]). \square

10. Classical coherence. In considering categories with operations and natural equivalences (replacing the equations of algebras) we immediately discover a relationship between connectedness and the commutativity of diagrams arising from the isomorphisms. This is illustrated by the following example.

Let \mathbf{V} be a category and $\otimes : \mathbf{V} \times \mathbf{V} \rightarrow \mathbf{V}$ a functor associative up to a natural isomorphism

$$a : A \otimes (B \otimes C) \rightarrow (A \otimes B) \otimes C$$

Since we have isomorphisms and not (in general) equalities it is then natural to ask whether all diagrams built up from the natural isomorphism a commute. This is an example of a *coherence* question. It turns out that the answer is affirmative in this case if a certain type of *pentagon* diagram commutes in a subcategory \mathbf{C} of the category of shapes $\mathcal{N}()$ that we are going to define next.

Shapes are defined inductively by

S1. 1 is a shape

S2. If T and S are shapes, so is $T \otimes S$

For each shape there is a variable set $\nu(T)$ defined inductively by

V1. $\nu(1)$ is a chosen one element set.

V2. $\nu(T \otimes S)$ is the disjoint union of $\nu(T)$ and $\nu(S)$.

Given ν , then for each shape T there is a functor $|T|$ given inductively by

F1. $|1| : \rightarrow$ is the identity functor.

F2. $|T \otimes S| = \otimes \cdot (|T| \times |S|) : W_T \times W_S \rightarrow \times \rightarrow$ for W_T and W_S products of $\nu(T)$ and $\nu(S)$ copies of ν , respectively.

Let $\mathcal{N}()$ be the category whose objects are all shapes and whose morphisms $F : T \rightarrow S$ are the natural transformations $F : |T| \rightarrow |S|$.

Given T, S and R we obtain $\alpha_{TSR} : T \otimes (S \otimes R) \rightarrow (T \otimes S) \otimes R$ in $\mathcal{N}()$ by letting $\alpha_{TSR}(X, Y, Z)$ be the component

$$|T|X \otimes (|S|Y \otimes |R|Z) \rightarrow (|T|X \otimes |S|Y) \otimes |R|Z$$

of the natural transformation a on \otimes for X, Y, Z objects of the domain categories for $|T|$, $|S|$ and $|R|$, respectively.

Let $\mathbf{C} (= \mathbf{C}())$ be the subcategory of $\mathcal{N}()$ whose objects are all shapes and whose morphisms, called the *allowable* morphisms of $\mathcal{N}()$, are given by

AM1. $1 : T \rightarrow T$ and $\alpha : T \otimes (S \otimes R) \rightarrow (T \otimes S) \otimes R$ are in \mathbf{C} for any shapes T, S, R .

AM2. If $f : T \rightarrow T'$ and $g : S \rightarrow S'$ are in \mathbf{C} , then so is $f \otimes g : T \otimes S \rightarrow T' \otimes S'$.

AM3. If $f : T \rightarrow S$ and $g : S \rightarrow R$ are in \mathbf{C} , then so is $gf : T \rightarrow R$.

Lemma. *Let \mathbf{D} be a category with a rank functor and assume that $(X/\mathbf{D})_{\mathcal{P}}$ is connected for each object X of \mathbf{D} . Then \mathbf{D} is a preorder if and only if the morphisms of \mathbf{D} are all monomorphisms.*

Proof. Let \mathbf{G} be the subcategory of all reduced objects. Then, since \mathbf{D} has a rank functor, every object of \mathbf{D} is \mathbf{G} -reducible. Thus, if $f, g : X \rightarrow Y$ in \mathbf{D} , then there is $h : Y \rightarrow G$ with G in \mathbf{G} . By Theorem 3.2 the Diamond Principle for \mathbf{G} holds. Thus objects of X/\mathbf{D} with codomain in \mathbf{G} are terminal. Thus $h \cdot f = h \cdot g$ since G is reduced and $f = g$ since h is monic. \square

A rank functor ρ for the subcategory \mathbf{C} of $\mathcal{N}()$ described above is defined recursively by $\rho(1) = 0$ and $\rho(T \otimes S) = \rho(T) + \rho(S) + |\nu(S)| - 1$, where $|\nu(S)| = \text{card } \nu(S)$.

Theorem. *The category $(X/\mathbf{C})_{\mathcal{P}}$ is connected for every shape $X \in |\mathbf{C}|$ provided that the pentagon diagram*

$$\begin{array}{ccccc}
 A \otimes (B \otimes (C \otimes D)) & \xrightarrow{a} & (A \otimes B) \otimes (C \otimes D) & \xrightarrow{a} & ((A \otimes B) \otimes C) \otimes D \\
 \downarrow 1 \otimes a & & & & \uparrow a \otimes 1 \\
 A \otimes ((B \otimes C) \otimes D) & \xrightarrow{a} & & \xrightarrow{a} & (A \otimes (B \otimes C)) \otimes D
 \end{array}$$

commutes for all A, B, C, D in \mathbf{C} .

Corollary. *The category \mathbf{C} is a preorder provided all pentagon diagrams commute.*

A functor $T : \mathbf{C} \rightarrow \mathbf{D}$ whose domain is a preorder is called a *coherence functor* for \mathbf{D} . Intuitively, T describes a class of commutative diagrams in \mathbf{D} .

The subcategory of $\mathcal{N}()$ generated by \mathbf{C} and the inverses to the associativity isomorphisms is also a preorder. For further details see ([6]) or ([8]).

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