

Math 342 Problem set 8 (due 11/3/09)

Prime rings

- Let R be a ring. We define a map $f: \mathbb{N} \rightarrow R$ inductively by $f(0) = 0_R$ and $f(n+1) = f(n) + 1_R$.
 - Show that $f(1) = 1_R$. Show that $f(n+m) = f(n) + f(m)$ for all $n, m \in \mathbb{N}$.
Hint: Induction on m .
 - Show that f respects multiplication, that is for all $n, m \in \mathbb{N}$, $f(nm) = f(n) \cdot f(m)$.
Hint: Induction again. The case $m = 0$ uses a result from class.OPTIONAL Extend f to a function $g: \mathbb{Z} \rightarrow R$ by setting $g(n) = f(n)$ if $n \in \mathbb{Z}_{\geq 0}$, and $g(n) = -f(-n)$ if $n \in \mathbb{Z}_{\leq 0}$. Show that g is a ring homomorphism.
Hint: Divide into cases.
- Let A, B be rings and $g: A \rightarrow B$ be a homomorphism. Show that the image $g(A) = \{b \in B \mid \exists a \in A : g(a) = b\}$ is a subring of B .
- Continuing problem 1, let g be the ring homomorphism you constructed, let $S = g(\mathbb{Z})$ be the image of g , and let $I = g^{-1}(0_R)$ be the set of $n \in \mathbb{Z}$ such that $g(n) = 0_R$.
 - Show that I is an ideal in \mathbb{Z} . By a previous problem set there is $m \in \mathbb{N}$ such that $I = (m)$.
 - If $m = 0$ show that g is injective, hence that R contains a subring isomorphic to \mathbb{Z} .
Hint: Use the criterion for injectivity from problem set 7.
 - Show that $m = 1$ is impossible, as long as $0_R \neq 1_R$.
Hint: What is $g(1)$ if $m = 1$? Compare with problem 1(a).
 - If $m \geq 2$, define $h: \mathbb{Z}/m\mathbb{Z} \rightarrow R$ by $h([a]_m) = g(a)$. Show that h is a well-defined function (that is, if $[a]_m = [a']_m$ then $g(a) = g(a')$).
 - Show that h is a ring homomorphism.
 - Show that h is an isomorphism.
Hint: To check injectivity, it is enough to understand $h([0]_m)$; to check surjectivity, given $s \in S$ need to find $[a]_m \in \mathbb{Z}/m\mathbb{Z}$ such that $h([a]_m) = s$.
We conclude that every ring contains either a subring isomorphic to \mathbb{Z} or a subring isomorphic to $\mathbb{Z}/m\mathbb{Z}$ for some $m \geq 2$.

REMARK. You can also check that $S = g(\mathbb{Z})$ is the smallest subring of R – the intersection of all subrings of R .

Prime fields and vector spaces

Now let F be a field, and let $g: \mathbb{Z} \rightarrow F$ be the map constructed in problem 1. Let m be the number defined in problem 3.

- Assume by contradiction that m is positive and composite, that is $m = ab$ with $1 < a, b < m$. Apply the function g and obtain a contradiction to the fact that F is a field. Conclude that either $m = 0$ or m is prime.

DEFINITION. m is called the *characteristic* of the field F and denoted $\text{char}(F)$. Problems 1-4 now show that the characteristic of a field is either zero or a prime number, and that a field of prime characteristic p contains an isomorphic copy of \mathbb{F}_p .

5. Let F be a finite field.
- Show that $\text{char}(F) > 0$. Conclude that $\mathbb{F}_p \subset F$ for some p .
Hint: You need to rule out $\text{char}(F) = 0$; for this use problem 3(b).
 - Show that F has the structure of a vector space over \mathbb{F}_p .
Hint: All the vector space axioms follow directly from the field axioms.
 - Show that $\dim_{\mathbb{F}_p} F < \infty$ (can F contain an infinite linearly independent set?). It follows that, as an \mathbb{F}_p -vector space, F is isomorphic to \mathbb{F}_p^n for some $n \geq 1$.
 - Show that the number of elements of a finite field is always a prime power.
Hint: How many elements are there in \mathbb{F}_p^n ?

REMARK. It is also true that for every $q = p^n$ there exists a field \mathbb{F}_q of size q , unique up to isomorphism.

The Hamming Code (variant)

6. (§13E.E6) Let $H \in M_{3 \times 7}(\mathbb{F}_2)$ be the matrix whose columns are all non-zero vectors in \mathbb{F}_2^3 , that is

$$H = \begin{pmatrix} 1 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{pmatrix}.$$

- Let $a, b, c, d \in \mathbb{F}_2$ be a 4-bit “message” we want to transmit. Show that there exist unique $x, y, z \in \mathbb{F}_2$ so that $H \cdot (x, y, z, a, b, c, d)^T = \underline{0}$. We will transmit the redundant 7-bit vector instead.
Hint: Need to show both that x, y, z exist and that they are unique.
- For each $1 \leq i \leq 7$, let \underline{e}^i be the standard basis vector of \mathbb{F}_2^7 with 1 at the i th co-ordinate. Calculate the seven vectors $H\underline{e}^i$.
- Let $\underline{v}, \underline{v}' \in \mathbb{F}_2^7$ be at Hamming distance 1. Show that there exists i so that $\underline{v}' = \underline{v} + \underline{e}^i$.
- Now let's say we transmit the 7-bit vector $\underline{v} = (x, y, z, a, b, c, d)^T$ from part (a) through a channel that can change at most one bit in every seven. Denote by \underline{v}' the 7 received bits, and show that if $\underline{v}' \neq \underline{v}$ then $H\underline{v}' \neq \underline{0}$. Conclude that the recipient can detect if a 1-bit error occurred.
Hint: Use the fact that $H\underline{v} = \underline{0}$ and your answers to parts (c) and (b).
- In fact, if at most one bit error can occur then the recipient can *correct* the error. Using the fact that the vectors $H\underline{e}^i$ are all different (see your answer to part (b)), show that knowing only \underline{v}' and that at most one error occurred, the recipient can calculate the difference $\underline{e} = \underline{v}' - \underline{v}$ and hence the original vector \underline{v} .
Hint: What are the possibilities for \underline{e} ? For $H\underline{e}$? how do they match up? Don't forget that it's possible that $\underline{v}' = \underline{v}$.