

3/12/2015

HW: half scored in [27, 40] / 60

Least time: Solvable groups:

$G$  solvable  $\Leftrightarrow$  normal series  $\{e\} \triangleleft G_0 \triangleleft G_1 \triangleleft \dots \triangleleft G_n = G$   
with  $G_{i+1}/G_i$  abelian.

Saws ①  $G$  solvable  $\Rightarrow H < G$ ,  $G/N$  solvable

②  $N \triangleleft G$ , if  $N$ ,  $G/N$  solvable so is  $G$ .

Example:  $B_n = \left\{ \begin{pmatrix} * & & * \\ & \ddots & \\ 0 & & * \end{pmatrix} \right\} \subset GL_n(F)$  (upper-triangular matrices)

HW:  $U_n \triangleleft B_n$ ,  $B_n/U_n \cong (F^*)^n$

Solvability top-down: relabel series  $G = G_0 \triangleright G_1 \triangleright G_2 \dots \triangleright G_n = \{e\}$

want  $G_0/G_1$  to be abelian, want for any  $x, y \in G$  that  $[\bar{x}, \bar{y}] = e$

$\bar{x} = xG_1$ ,  $\bar{y} = yG_1$  images of  $x, y$  in  $G/G_1$ .

$\Leftrightarrow$  want  $[x, y] \in G_1$ .

$\Rightarrow G/G_1$  abelian iff  $G_1 \supset \{[x, y] \mid x, y \in G\}$

iff  $G_1 \supset \langle [x, y] \rangle \stackrel{\text{def}}{=} [G, G] = G'$

HW: If  $G = G_0 \triangleright G_1 \triangleright \dots \triangleright G_n = \{e\}$   $\uparrow$  "derived subgroup"

$G_i/G_{i+1}$  abelian then  $G_i \supseteq G^{(i)}$ :  $G^{(0)} = G$ ,  $G^{(i+1)} = (G^{(i)})'$

clearly  $G^{(i)}/G^{(i+1)}$  commutative: killed all commutators

conclusion:  $G$  solvable iff  $G^{(n)} = \{e\}$  for some  $n$ .

$G^{(1)}$  generated by  $[x, y]$

$G^{(2)}$  " "  $[x, y], [z, w]$

checks  $(B_n)' = U_n$

corollary:  ~~$SL_n(\mathbb{F})$~~   $(GL_n(\mathbb{F}))' \supseteq (B_n)' = U_n$

also:  $\det(xy x^{-1} y^{-1}) = 1$ , so  $[x, y] \in \ker(\det) = SL_n(\mathbb{F}) = \{g \in GL_n(\mathbb{F}) \mid \det g = 1\}$

$G'$  normal;  $\varphi \in \text{Aut}(G)$ . Then  $\varphi([x, y]) = \varphi(xy x^{-1} y^{-1})$   
 $= \varphi(x) \varphi(y) \varphi(x^{-1}) \varphi(y^{-1})$   
 $= [\varphi(x), \varphi(y)]$

then  $\varphi$  permutes the commutators,

so fixed the subgp they generate:  $\varphi(G') = G'$

What is the normal subgp generated by  $U_n(\mathbb{F})$ ?

over  $\mathbb{R}$ ,  $U_n(\mathbb{R}) =$  upper triangular +1-diagonal,  $\bar{U}_n =$  lower-triangular +1-diagonal

$$\begin{pmatrix} 1 & * & * \\ & 1 & * \\ & & 1 \end{pmatrix}$$

$$\begin{pmatrix} 1 & & \\ * & 1 & \\ * & * & 1 \end{pmatrix}$$

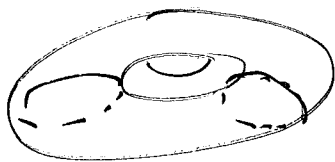
jointly generate  $SL_n(\mathbb{R})$  (Gaussian elimination)

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# Group theory in topology

Topology: study of properties unchanged by deformation

Basic questions: given  $X, Y$  want to know: are they  
distinguish  $X, Y$  say, by dimension the same!



Let  $\pi_1$  be fundamental group

$(X, *)$  is a (topological space) + pt  $*$

consider the set of based loops  $C((S^1, *), (X, *))$

(cts maps  $\gamma: [0, 1] \rightarrow X$ , s.t.  $\gamma(0) = *, \gamma(1) = *$ )

natural operation: concatenation: if  $\gamma_1, \gamma_2$  loops define

$\gamma_1 \cdot \gamma_2$  : do  $\gamma_1$ , then  $\gamma_2$ .

associative:  $(\gamma_1 \cdot \gamma_2) \cdot \gamma_3 = \gamma_1 \cdot (\gamma_2 \cdot \gamma_3)$

identity: constant loop  $e(t) = *$  for all  $t$ .

inverse:  $\bar{\gamma}(t) = \gamma(1-t)$  (reverse direction)

declare  $\gamma_1 \sim \gamma_2$  if can deform  $\gamma_1$  to  $\gamma_2$  (endpoints fixed)

deform  $\gamma \cdot \bar{\gamma}$  to  $e$ : at time 0



at time  $\epsilon$



at time  $1-\epsilon$



at time 1



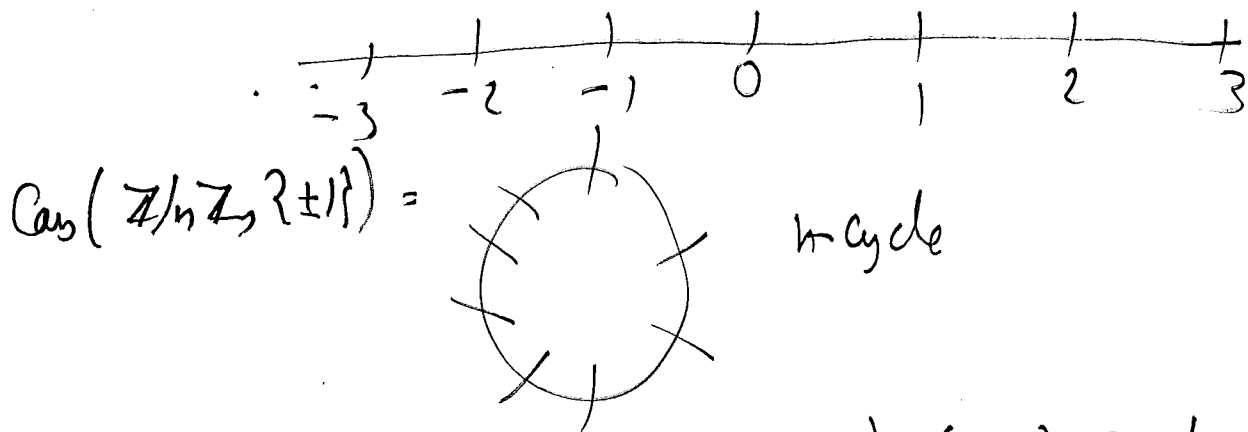
# Combinatorics

$G$  gp,  $S$  generating set. (assume symmetric: if  $s \in S$ ,  $s^{-1} \in S$  too)

Def: the Cayley graph  $\text{Cay}(G; S)$  is the graph with vertex set  $G$   
edges  $E = \{(g, gs) \mid \begin{matrix} g \in G \\ s \in S \end{matrix}\}$

(Graph: pair  $\Gamma = (V, E)$   $V$  vertices,  $E$ : edge connect vertices)

Example  $\text{Cay}(\mathbb{Z}, \{\pm 1\})$  edges:  $(n, n \pm 1)$



if  $(g, gs) \in E$  then  $(gs, (gs)s^{-1}) = (gs, g) \in E$  too

$G$  acts on  $\text{Cay}(G; S)$ : translation  $g \cdot x = gx$   
maps edge  $(x, xs)$  to edge  $(gx, (gx)s)$

Example: (subject to a congruence condition)  $SL_2(\mathbb{F}_q)$ ,  $q$  prime,  
has a generating set  $S$  of size  $p+1$  with  $\text{Cay}(SL_2(\mathbb{F}_q); S)$   
extremely well-connected:

This network has  $\approx q^3$  vertices, each with only  $p+1$  neighbours  
(think  $p$  fixed,  $q \rightarrow \infty$ ) but if cut it into two pieces  $A, B$   
many connections across the cut.

check ①  $\sim$  is an equiv rel

② If  $\gamma_1 \sim \gamma'_1$   
 $\gamma_2 \sim \gamma'_2$  then  $\gamma_1 \cdot \gamma_2 \sim \gamma'_1 \cdot \gamma'_2$

Follows: • still well-defined on  $C((S', x), (X, y)) / \sim$   
[e] still identity

$$\text{now } [\gamma] \cdot [\bar{\gamma}] = [\gamma \cdot \bar{\gamma}] = [e]$$

so we got a group! Call it  $\pi_1(X, *)$ .

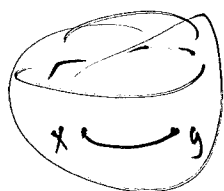
Examples:  $\pi_1(S^1) = \mathbb{Z}$  (a loop on  $S^1$ , need to know only winding number)

$$\pi_1(S^2) = \{e\}$$

$$\pi_1(S^1 \times S^1) = \mathbb{Z} \times \mathbb{Z} = \mathbb{Z}^2$$

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suppose  $x, y$  points in  $X$  connected by path  $p$



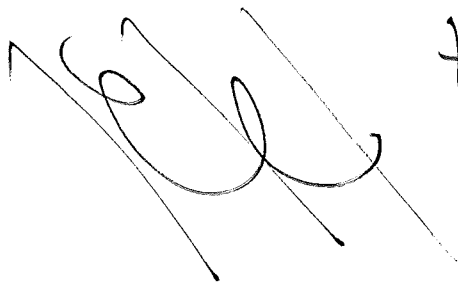
given loop  $\gamma$  based at  $y$   
the loop  $p \cdot \gamma \cdot \bar{p}$  is based at  $x$

converse identification uses  $\bar{p}$ ,  
these are inverse modulo deformation:

$$\pi_1(X, x) \cong \pi_1(X, y)$$

$$\pi_1(X, x) = (C((S', x), (X, x)) / \sim, \cdot)$$

$$\pi_1(X \times Y) = \pi_1(X) \times \pi_1(Y)$$



$$\mathbb{H} = \{z = x + iy \mid y > 0\}$$

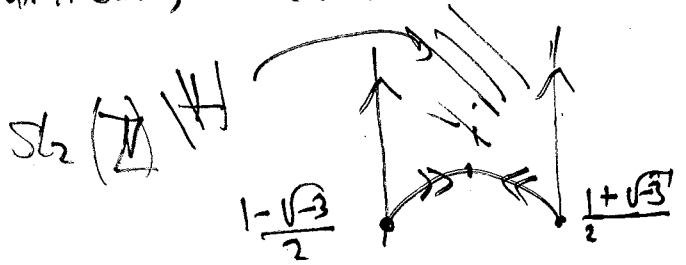
$g \in SL_2(\mathbb{R}) : g = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$  define  $g \cdot z = \frac{az + b}{cz + d}$

ex: this is a gp action on  $\mathbb{H}$ .

(preserve distance) action transitive,

$$\text{Stab}_{SL_2(\mathbb{R})}(i) = SO(2)$$

In particular,  $SL_2(\mathbb{Z})$  acts.



$$\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} z = z + 1$$

$$\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} z = -\frac{1}{z}$$

example count solutions to  $x^2 + y^2 + z^2 + w^2 = N$  ( $(x, y, z, w) \in \mathbb{Z}^4$ )  
 using holomorphic fns on  $SL_2(\mathbb{Z}) \backslash \mathbb{H}$